

TALK 3 : Section 2.3 of FFRS 2006

(by H. Yadav)

Last time : Defn. of MTC  
Idempotents and retracts, Karoubi completion of  $\mathcal{C}$

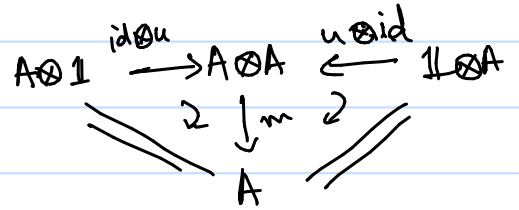
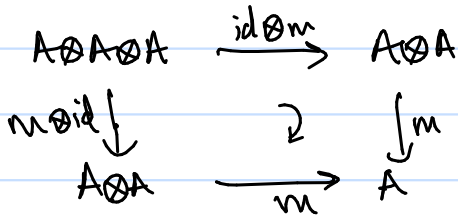
GOAL : This section defines the different algebraic structures that will be used throughout the paper.

(basically we define the usual algebraic notions like algebra, modules in monoidal categories.

Let  $(\mathcal{C}, \otimes, \mathbb{1})$  be a strict monoidal category.

Defn : An algebra in  $\mathcal{C}$  is a triple  $(A, m, u)$  where

- $A \in \text{obj}(\mathcal{C})$
- $m : A \otimes A \rightarrow A, u : \mathbb{1} \rightarrow A \in \text{Mor}(\mathcal{C})$
- $m$  is associative  $u$  is unital



commute

Writing commutative diagrams, <sup>takes time</sup> so we use graphical calculus (it is also very visual!)

Rules:  $X \xrightarrow{f} Y$   $\begin{array}{c} \textcircled{f} \\ | \\ X \end{array}^Y$ ,  $X \xrightarrow{f} Y, W \xrightarrow{g} Z$  then  $f \otimes g = \begin{array}{c} Y \quad Z \\ \textcircled{f} \quad \textcircled{g} \\ | \quad | \\ X \quad W \end{array}$

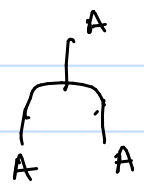
$Y \xrightarrow{h} Z$ , then  $h \circ f = \begin{array}{c} Z \\ \textcircled{h} \\ \textcircled{f} \\ | \\ X \end{array}$ , don't usually draw the unit object

unit map of algebra

$$\mathbb{1} \xrightarrow{\eta} A$$

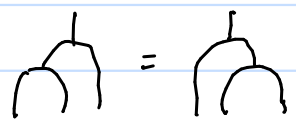


multiplication

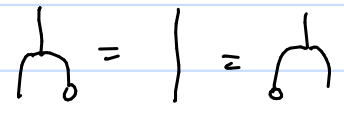


Using graphical notation,

associativity is



unitality is



Defn: A coalgebra in  $\mathcal{C}$  is a triple  $(C, \Delta, \epsilon)$  where

- $C \in \text{Obj}(\mathcal{C})$
- $\Delta: C \rightarrow C \otimes C$  ( $\begin{array}{c} C \\ \cup \\ C \otimes C \end{array}$ ),  $\epsilon: C \rightarrow \mathbb{1}$  ( $\begin{array}{c} C \\ \downarrow \\ \mathbb{1} \end{array}$ )  $\in \text{Mor}(\mathcal{C})$
- $\Delta, \epsilon$  satisfy

Defn: A left module over an algebra  $A$  in  $\mathcal{C}$  is a pair  $M = (M, \rho)$  where

- $M \in \text{Obj}(\mathcal{C})$
- $\rho: A \otimes M \rightarrow M$  ( $\begin{array}{c} M \\ \downarrow \\ A \otimes M \end{array}$ ) is a morphism satisfying

Defn: Take two left  $A$ -modules  $(M, \rho_M), (N, \rho_N)$ . Then  $f: M \rightarrow N$  is called a map of left  $A$ -modules if  $f$  satisfies

left  $A$ -modules and  $A$ -module maps form a category which we denote as  $\mathcal{C}_A$ .

→ Can similarly define right  $A$ -modules

- $M \in \mathcal{C}_A$  is called simple if it simple as an object of  $\mathcal{C}_A$
- For  $U \in \text{ob}(\mathcal{C})$ , denote  $\text{Ind}_A(U) = A \otimes U$ . Then  $\text{Ind}_A(U) \in \mathcal{C}_A$  via the multiplication in  $A$ . Denote by  $\mathcal{C}_A^{\text{ind}} \subset \mathcal{C}_A$ , the full subcategory of  $\mathcal{C}_A$  whose objects are  $\text{Ind}_A(U)$  for some  $U \in \text{ob}(\mathcal{C})$

Defn: An  $A$ -bimodule is a triple  $M = (M, \rho_L, \rho_R)$  such that  $(M, \rho_L)$  is a left  $A$ -module and  $(M, \rho_R)$  is a right  $A$ -module, and left and right actions commute.

$$\rho_L = \begin{array}{c} M \\ \swarrow \downarrow \searrow \\ A \quad M \end{array}, \quad \rho_R = \begin{array}{c} M \\ \swarrow \downarrow \searrow \\ M \quad A \end{array}$$

then

$$\begin{array}{c} M \\ \swarrow \downarrow \searrow \\ A \quad M \quad A \end{array} = \begin{array}{c} M \\ \swarrow \downarrow \searrow \\ A \quad M \quad A \end{array}$$

→  $A$ -bimodule and bimodule maps (i.e. left & right module maps) between them form a category that we will denote as  $\mathcal{C}_{A|A}$ .

→  $\mathcal{C}_{A|A}$  is monoidal with

$$\left( \begin{array}{c} M \\ \swarrow \downarrow \searrow \\ A \quad M \end{array}, \begin{array}{c} N \\ \swarrow \downarrow \searrow \\ M \quad N \end{array} \right) \otimes \left( \begin{array}{c} N \\ \swarrow \downarrow \searrow \\ N \quad N \end{array}, \begin{array}{c} N \\ \swarrow \downarrow \searrow \\ N \quad N \end{array} \right) = \left( \begin{array}{c} M \quad N \\ \swarrow \downarrow \searrow \\ A \quad M \quad N \end{array}, M \otimes N, \begin{array}{c} M \quad N \\ \swarrow \downarrow \searrow \\ M \quad N \quad A \end{array} \right)$$

Defn A Frobenius algebra in  $\mathcal{C}$  is a 5-tuple  $(A, m, u, \Delta, \varepsilon)$  such that

- $(A, m, u)$  is an algebra in  $\mathcal{C}$
- $(A, \Delta, \varepsilon)$  is a coalgebra in  $\mathcal{C}$
- $\Delta \circ m$  satisfy the Frobenius law

$$\cup = \cap = \cup$$

- A Frobenius algebra is called special if  $\varepsilon \circ \eta = \beta_{\mathbb{1}} \text{id}_{\mathbb{1}}$ ,  $m \circ \Delta = \beta_A \text{id}_A$

$$\left( \int = \beta_{\mathbb{1}} \text{id}_{\mathbb{1}} \quad \oint = \beta_A \text{id}_A \right)$$

for some non-zero numbers  $\beta_{\mathbb{1}}$  and  $\beta_A$ .

### Updating graphical calculus for $\mathcal{C}$ -rigid

Recall that if  $\mathcal{C}$  is rigid, then  $\forall U \in \mathcal{C}$ , we have

$$\left( U^v, b_u = \overset{v}{\cup} \overset{u}{\cup}, d_u = \underset{u}{\cap} \underset{v}{\cap} \right) \text{ and } \left( {}^v U, \tilde{b}_u = \overset{v}{\cup} \overset{u}{\cup}, \tilde{d}_u = \underset{u}{\cap} \underset{v}{\cap} \right)$$

RIGHT DUAL LEFT DUAL

satisfying

$$\overset{u}{\cup} \underset{u}{\cap} = |_u, \underset{u^v}{\cap} \overset{u^v}{\cup} = |_{u^v} \text{ and } \overset{v}{\cup} \underset{v}{\cap} = |_v, \underset{v^u}{\cap} \overset{v^u}{\cup} = |_{v^u}$$

Defn: (i) An algebra  $A$  in rigid monoidal category  $\mathcal{C}$  together with a morphism  $\varepsilon: A \rightarrow \mathbb{1}$  (denoted  $\int$ ) is called a symmetric algebra if

$$\Phi_1 := \int \overset{A^v}{\cup} \underset{A}{\cap} = \int \overset{A}{\cap} \underset{A^v}{\cup} =: \Phi_2 \quad (*)$$

(ii) A special symmetric Frobenius algebra is a special Frobenius algebra  $A$  that is symmetric with  $\varepsilon = \text{counit of } A$ .

→ For a special symmetric Frobenius algebra  $A$ , we have  $\dim(A) = \int = \oint = \beta_A \int = \beta_A \beta_{\mathbb{1}}$

→ We normalize  $\epsilon, \Delta$  so that  $\beta_{11} = \dim(A)$ ,  $\beta_A = 1$   
 This convention is followed unless mentioned otherwise.

## Various remarks about Frobenius algebras:

•  $\mathcal{C}$  semisimple,  $A$  special Frobenius  $\Rightarrow \mathcal{C}_A$  is semisimple

• We have the induction & restriction functors

$$\begin{array}{ccc} \text{Ind: } \mathcal{C} & \longrightarrow & \mathcal{C}_A \\ U & \longmapsto & (A \otimes U, \Delta|) \end{array} \qquad \begin{array}{ccc} \text{Res: } \mathcal{C}_A & \longrightarrow & \mathcal{C} \\ M = (M, \rho) & \longmapsto & M \end{array}$$

For any algebra  $A$ , we have reciprocity relation that is, a canonical natural bijection

$$\text{Hom}_A(\text{Ind}_A(U), M) \xrightarrow{\cong} \text{Hom}(U, M)$$

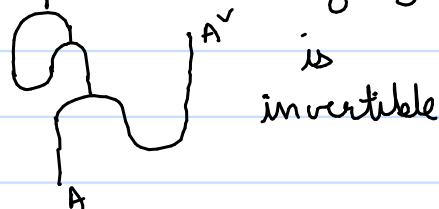
Categorically, this amounts to saying that

$$\text{Ind}_A \dashv \text{Res} \quad (\text{left adjoint})$$

If  $A$  is Frobenius,  $\text{Res} \dashv \text{Ind}$  as well.

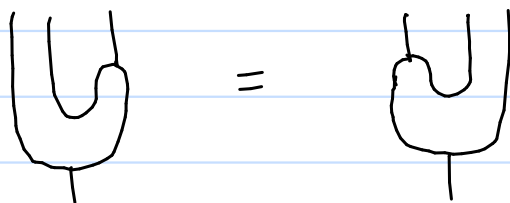
• Given an algebra  $(A, m, u)$  in a pivotal category,

$\exists \Delta, \epsilon$  making  $(A, m, u, \Delta, \epsilon)$  symmetric special Frobenius  $\Leftrightarrow$



• In terms of graphical calculus, an algebra being symmetric Frobenius means that multiplication and comultiplication can be moved past each other in all possible arrangements.

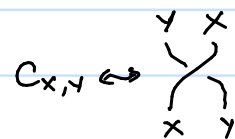
example



Now  $\mathcal{C}$  is braided as well.

Then we have a natural isomorphism

$$\{C_{X,Y} : X \otimes Y \rightarrow Y \otimes X\}_{X,Y \in \text{ob}(\mathcal{C})}$$



Defn: For  $A$  an algebra in a braided tensor category  $\mathcal{C}$ , we have the functors

$$\alpha_A^\pm : \mathcal{C} \longrightarrow \mathcal{C}_{A|A}$$

of  $\alpha$ -induction.

$$\alpha_A^+(V) := (A \otimes V, \text{left action}, \text{right action}) ; \alpha_A^+(f) = \text{id}_A \otimes f$$

$$\alpha_A^-(V) := (A \otimes V, \text{left action}, \text{right action}) ; \alpha_A^-(f) = \text{id}_A \otimes f$$

Remark:  $\alpha^\pm : \mathcal{C} \rightarrow \mathcal{C}_{A|A}$  are monoidal functors.

Defn: An algebra  $A$  in a braided monoidal category is said to be commutative if

$$m \circ C_{A,A} = m$$

Remark:  $\rightarrow$  For general  $A$ ,  $\mathcal{C}_A$  is not monoidal.

$\rightarrow$  But for  $A$  commutative,  $\mathcal{C}_A$  is monoidal. There are two choices for monoidal structure.

$\rightarrow$  Later study  $\mathcal{C}_A^{\text{loc}}$ , full subcategory of  $\mathcal{C}_A$ . It has a canonical tensor structure.

Now say  $\mathcal{C}$  is a braided pivotal category. Then the twist  $\theta$  satisfies

$$\theta = \text{diagram of a twist on a strand labeled A}$$

PROP 2.25 Let  $\mathcal{C}$  be a braided pivotal category

(i) A commutative symmetric Frobenius algebra has trivial twist, i.e.,  $\Theta_A = \text{id}_A$ .

(ii) Every commutative Frobenius algebra with trivial twist is symmetric.

(iii) A commutative symmetric Frobenius algebra is also cocommutative.

Proof of (i):  $A$  is symmetric

$$\Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array} \quad \begin{array}{l} \text{apply } d_A = \begin{array}{c} \text{A} \\ \text{A} \end{array} \\ \Rightarrow \\ \text{to both sides} \end{array} \quad \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{using snake} \\ \text{relation on} \\ \text{RHS} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{moving the} \\ \text{left thread over} \\ \text{mult. map} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{move the shaded} \\ \text{part to left side} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{Commutativity} \\ \text{of } A \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{adding extra} \\ \text{maps to both} \\ \text{sides} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{Frobenius} \\ \text{law} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

$$\left( \begin{array}{l} \text{by Unit / Counit} \\ \text{axioms} \end{array} \right) \Rightarrow \begin{array}{c} \text{A} \\ \text{A} \end{array} = \begin{array}{c} \text{A} \\ \text{A} \end{array}$$

See paper for (ii), (iii)

