

proof of alternating linear clobber conj.

hayward @ ualberta . ca

XChen TFolkersen KHasham RH DLee

ORandall LSchultz EVandermeer

Bonn Forschungsinst-Diskrete-Math

Oberseminar Sep 15 2025

who wins chess?

who wins chess?

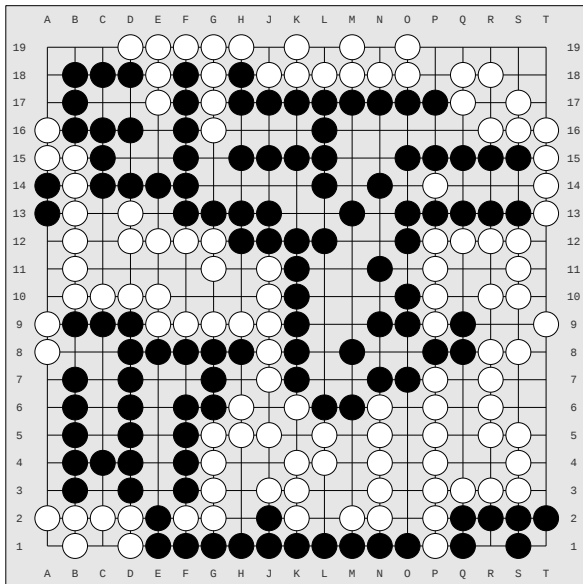
- maybe 1st player ?
- maybe draw ?
- can you prove either of these ?

who wins chess?

- Zermelo's theorem 1913:
- finite alt-turn win-lose-draw game, either
1st-player can always win, or
2nd-player can always win, or
each player can always at-least-draw.

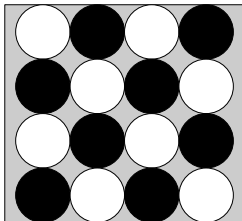
- Nim
- Sprague-Grundy
- watching Go end games, 1970s
- Berlekamp-Conway-Guy: Winning Ways
- Conway: On Numbers and Games

Berlekamp/Wolfe 9-dan stumper

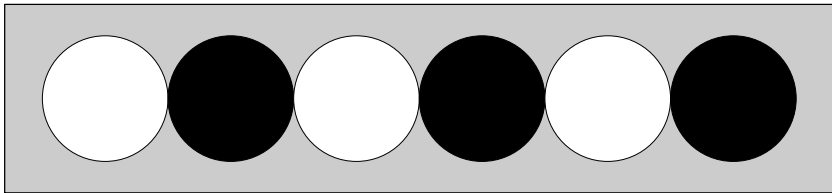


Mathematical Go:
Chilling Gets the Last Point
(Berlekamp/Wolfe)

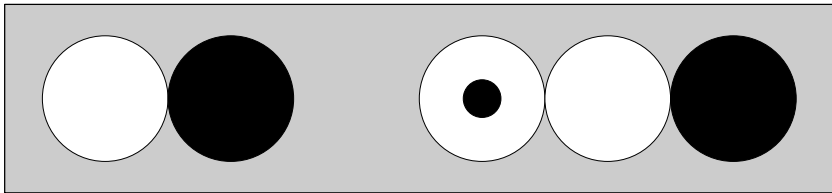
x plays first, who wins?



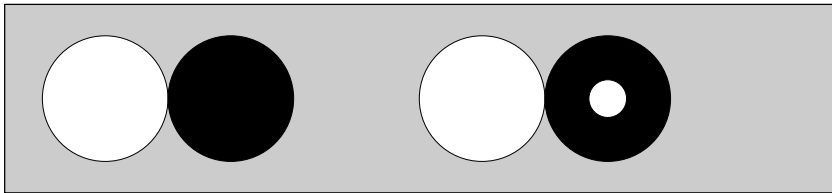
example Alternating Linear Clobber game



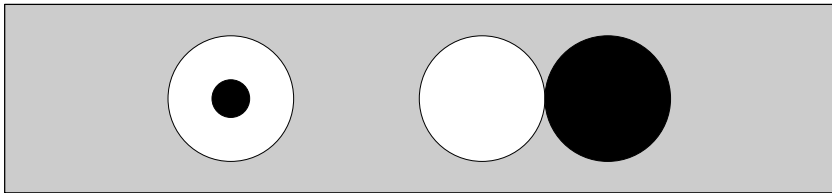
example Alternating Linear Clobber game



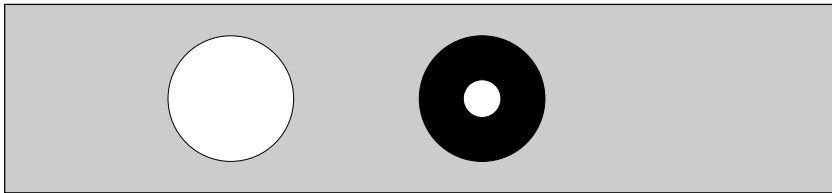
example Alternating Linear Clobber game



example Alternating Linear Clobber game

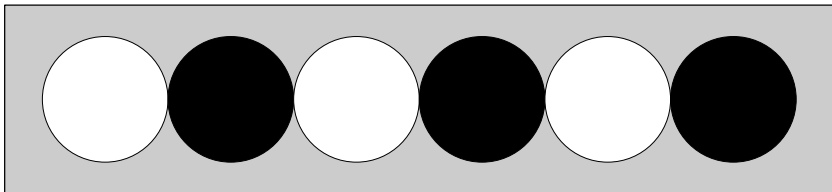


example Alternating Linear Clobber game



now White has no legal moves, White loses

play a game with your neighbour



- *intro to clobber* 2001 AGNW

Albert Grossman Nowakowski Wolfe

<https://webdocs.cs.ualberta.ca/~hayward/papers/AGNW.pdf>

this talk:

<https://webdocs.cs.ualberta.ca/~hayward/talks/cgtcv.pdf>

alternating linear clobber

- *linear clobber* clobber on a path
- *alternating linear clobber*

start from alternating sequence

OX, OXOX, OXOXOX, OXOXOXOX, ...

conjecture 3.2 AGNW 2004:

for each ALC start position except $oxoxox$,
first player can win

we prove this



how to prove this? a visual riddle

how to prove this? a visual riddle



The most environmentally friendly way to cook! More efficient than any other cooking technology

how to prove this? a visual riddle

proof

by

induction

combinatorial games: zero, neg, sum

- zero $\{ \mid \}$
- sum $g + h$
- negation
- equivalence $g \equiv h$ iff $g + -h$ is $\mathcal{P}$ -position

combinatorial games: outcome classes

- $\mathcal{P}$

- $\mathcal{L}$

- $\mathcal{R}$

- $\mathcal{N}$

combinatorial games: equality

- $g = h$ iff game trees isomorphic
- example
- $g = \{a, b, c \mid d, e\}$
- $h = \{b, a, c \mid e, e\}$

- players L, R , assume L plays 1st
- $\mathcal{A}$ set of starting ALC positions
- K some games reachable from $\mathcal{A}$
- for each k in K , for each R -move $k \rightarrow h$,
find L -move $h \rightarrow k'$ or $h \rightarrow 0$
- details: $K?$ L -win strategy?

6 sets of parts: $\mathcal{A}$, $\mathcal{O}$, $o\mathcal{A}$

$$\mathcal{A} = \{ \text{ox}, \text{oxox}, \text{oxoxox}, \dots \}$$

$$\{ \text{a2}, \text{a4}, \text{a6}, \dots \}$$

$$\mathcal{O} = \{ \text{o}, \text{oxo}, \text{oxoxo}, \dots \}$$

$$\{ \text{o}, \text{o3}, \text{o5}, \dots \}$$

$$o\mathcal{A} = \{ \text{o}, \text{oox}, \text{ooxox}, \dots \}$$

$$\{ \text{o}, \text{oo3}, \text{oo5}, \dots \}$$

6 sets of parts: $o\mathcal{O}$, $o\mathcal{O}o$, $o\mathcal{A}x$

$$o\mathcal{O} = \{oo, \quad ooXo, \quad ooXoXo, \quad \dots\}$$

$$\{oo, \quad oo4, \quad oo6, \quad \dots\}$$

$$o\mathcal{O}o = \{ooo, \quad ooXoo, \quad ooXoXoo, \quad \dots\}$$

$$\{ooo, \quad oo5oo, \quad oo7oo, \quad \dots\}$$

$$o\mathcal{A}x = \{ooxx, \quad ooXoxx, \quad ooXoXoxx, \quad \dots\}$$

$$\{ooxx, \quad oo6xx, \quad oo8xx, \quad \dots\}$$

$$\text{OX}, \text{OXO}, \text{OOXOX}, \text{OOXOXOO} = *$$

$$\text{OXX} = \uparrow$$

$$\text{OXOX} = \pm\{\uparrow, 0\}$$

$$\text{OOXO}, \text{OXOXOXOXO} = \uparrow*$$

$$\text{OOXOXX}, \text{OXOXOXOXOXOX} = \text{OXOX} *$$

$$\text{OXOXO} = \{\uparrow \mid *\}$$

standard form: apply after each move

o, oo, ooo, ooxx, oxoxox $\Rightarrow$ - (game 0)

oxo, ooxox, ooxoxoo $\Rightarrow$ ox

oxoxoxoxo $\Rightarrow$ ooxo

ooxoxx, oxoxoxoxoxox $\Rightarrow$ oxox ox

ooxoo $\Rightarrow$ oox

ooxo $\Rightarrow$ xxo ox

standard form

- also apply the negative rule forms

e.g. $xxox \Rightarrow oox \quad xo$

- also apply $p + -p \Rightarrow -$ (game 0)

e.g. $xxo + oox \Rightarrow -$

- repeatedly apply rules until no rule applies

theorem 1 (easy)

theorem 1:

all parts that arise in ALC game
in one of $\mathcal{A}$, $\mathcal{O}$, $\circ\mathcal{A}$, $\circ\mathcal{O}$, $\circ\mathcal{O}\circ$, $\circ\mathcal{A}x$
or their negatives.

$\mathcal{A}'$ to $o\mathcal{O}o'$

$$\mathcal{A}' = \mathcal{A} \setminus \{a_2, a_4, a_{12}\}$$

$$\mathcal{O}' = \mathcal{O} \setminus \{o, o_3, o_9\}$$

$$o\mathcal{O}' = o\mathcal{O} \setminus \{oo, oo_4, oo_6, oo_8\}$$

$$o\mathcal{A}' = o\mathcal{A} \setminus \{o, oo_3, oo_5\}$$

$$o\mathcal{O}o' = o\mathcal{O}o \setminus \{ooo, oo_7oo\}$$

$\mathcal{A}'$ to $o\mathcal{O}o'$

$\mathcal{A}'$ { a8 a10 a14 ... }

$\mathcal{O}'$ { o5 o7 o11 o13 ... }

$o\mathcal{O}'$ { oo10 oo12 oo14 ... }

$o\mathcal{A}'$ { oo7 oo9 oo11 oo13 oo15 ... }

$o\mathcal{O}o'$ { oo5oo oo9oo oo11oo oo13oo ... }

definition of irregular games $\mathcal{I}$

$$\begin{aligned}\mathcal{I} &= \{a_2, a_4, oo_6\} \\ &= \{ox, oxox, ooxoxo\}\end{aligned}$$

definition of set K

$K =$ games g , each part in any of

$\mathcal{O}' \quad o\mathcal{O}' \quad o\mathcal{O}o' \quad \mathcal{I} \quad \{xxo\} \quad \{oo8\} \quad \mathcal{A}' \quad o\mathcal{A}'$

(a, b, c, d, e, f, y, z)

count vector $\mu(g)$

definition of games to be avoided Q

- $Q = \{o_5 + a_2, o_7 + a_4, o_{15} + a_4 + a_2\}$
- thm: $\mathcal{O}' \subseteq \mathcal{L}$
- $\text{sum } \mathcal{O}' + ox \setminus Q \subseteq \mathcal{L}$
- $\text{sum } \mathcal{O}' + oxox \setminus Q \subseteq \mathcal{L}$
- $\text{sum } \mathcal{O}' + oxox + ox \setminus Q \subseteq \mathcal{L}$

definition of sets $\mathcal{S}_0, \mathcal{S}_1, \mathcal{S}_2$

$\mathcal{S}$ games of $K \setminus Q$

$\mathcal{S}_0$: $a \geq c, \quad y \leq 1, \quad z = 0 \quad \text{or}$

$a \geq c + 1, \quad y = 0, \quad z = 1$

$\mathcal{S}_1$: $y, z = 0, \quad a \geq c + 1$

$\mathcal{S}_2$: $a, b, c, y, z = 0, \quad e \geq 1$ and

$d = 0 \quad \text{or} \quad d + e \geq 3$

definition of set LL

- LL is set of three games

$\{ \text{oo8},$

$\text{o5} + \text{oox} + \text{a4} + \text{a2},$

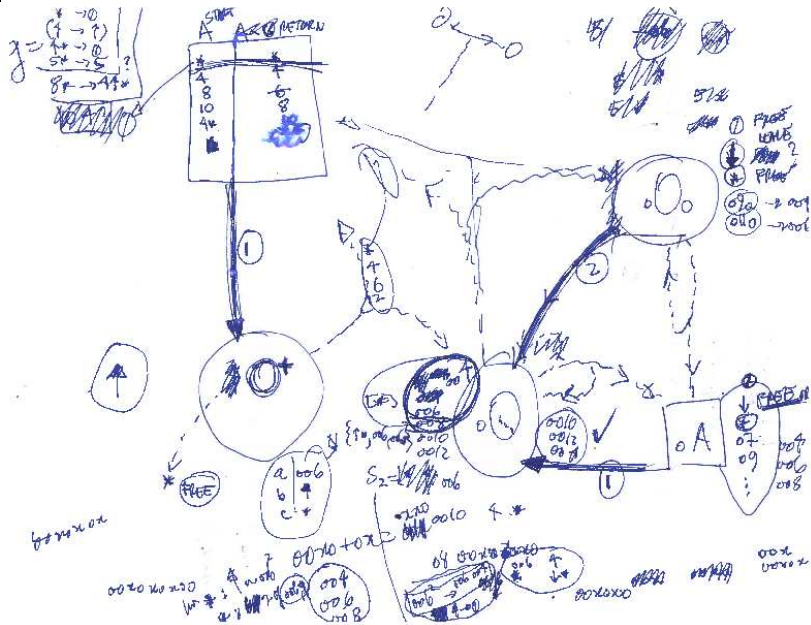
$\text{a14} + \text{xxo} + \text{a4} + \text{a2} \quad \}$

- $g \in \mathcal{S}_1 \cup \mathcal{S}_2$: R-move $\rightarrow \mathcal{S}_0$
- $g \in \mathcal{S}_0$: L-move $\rightarrow \mathcal{S}_1 \cup \mathcal{S}_2 \cup LL \cup \{0\}$
- $\mathcal{S}_0 \subset \mathcal{L} \cup \mathcal{N}$
- $(\mathcal{S}_1 \cup \mathcal{S}_2) \subset \mathcal{L}$
- $\mathcal{A} \setminus \{\text{oxoxox}\} \subset \mathcal{N}$

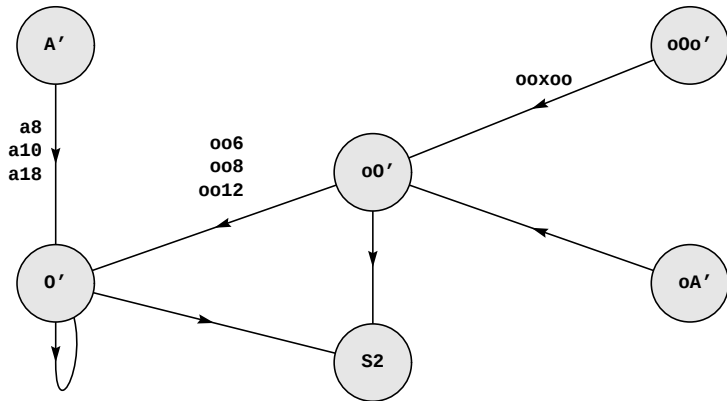
how L wins on $\mathcal{A} \setminus \{\text{oxoxox}\}$: outline

- play any part in $\mathcal{A}' \cup o\mathcal{O}'$ to $\mathcal{O}'$
- play any part in $o\mathcal{O}o' \cup o\mathcal{A}'$ to $o\mathcal{O}'$
- avoid Q
- ignore $\mathcal{I}$ until the end

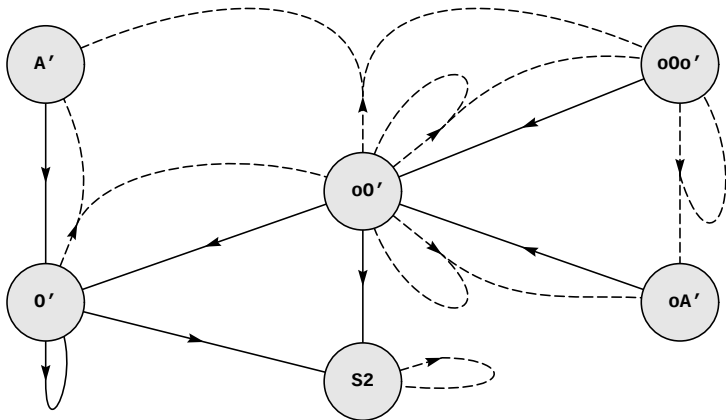
how L wins



how L wins



how L and White play



where is the devil?

where is the devil?

in the details

<https://arxiv.org/abs/2509.08985>

how L wins

rule	{game} or subgame	result
1	$\mathcal{A}'$ non-empty	
a	$\{a_8 + a_2\}$	$xxo + a_4 + a_2$ (avoid $o_5 + a_2$)
b	$\{a_{10} + a_4\}$	$o_5 + xxox + a_4$ (avoid $o_7 + a_4$)
c	$\{a_{18} + a_4 + a_2\}$	$a_{14} + xxo + a_4 + a_2$ (avoid $o_{15} + a_4 + a_2$)
d	any of $a_8, a_{10}, a_{14}, \dots$	$o_5, o_7, o_{11}, \dots$

how L wins

rule	{game} or subgame	result
2	$\mathcal{A}'$ empty, $o\mathcal{A}'$ non-empty	
	oo7, oo9, ...	oo4, oo6, ...

how L wins

rule	{game} or subgame	result
3	$\mathcal{A}'$, $o\mathcal{A}'$ empty, $o\mathcal{O}o'$ non-empty	
a	$\{o5 + oox\}$	$xxo + oox$
b	$o5 + a4 + a2 + oox$	$o5 + a4 + a2 + a2$
c	$a4 + oox$	$xxo + oox$
d	oox	$a2$
e	$oo9oo, oo11oo, \dots$	$oo4, oo6, \dots$

how L wins

rule	{game} or subgame	result
4	$\mathcal{A}', o\mathcal{A}', o\mathcal{O}o'$ empty, $o\mathcal{O}'$ non-empty	
a	{oo10 +a2}	o7 +a2 +a2 (avoid o5 +a2)
b	{oo12 +a4}	oo8 +xxo +a4 (avoid o7 +a4)
c	oo10	o5
d	oo12	o7
e	oo14	o11 +a2
f	oo16	o11
g	oo18	o13 (avoid o15 +a4 +a2)
h	oo20	o17 a2 (avoid o15 +a4 +a2)
i	oo22, oo24, ...	o17, o19, ...

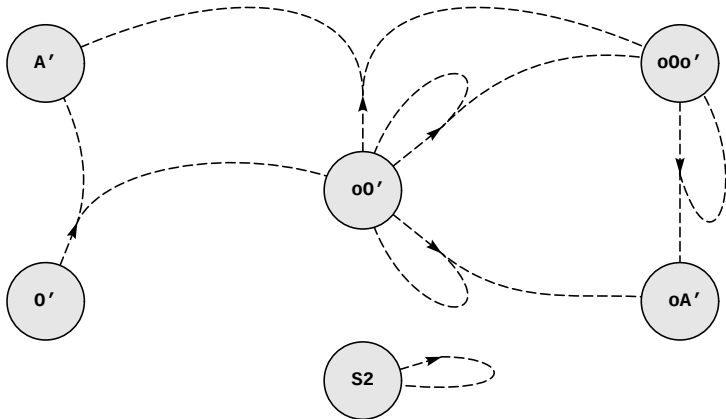
how L wins

rule	{game} or subgame	result
5	$\mathcal{A}', o\mathcal{A}', o\mathcal{O}o', o\mathcal{O}'$ empty, $\mathcal{I}$ non-empty	
a	$oo6 + a4$	$xxo + oo6 + a4$
b	$oo6 + oo6 + a2$	$xxo + oo6 + a2$
c	$oo6 + oo6$	$ooxo + oo6 = xxo + oo6 + a2$
d	$oo6 + a4 + a2$	$xxo + a4 + a2$
e	$oo6 + a4$	$ooxo + a4 = xxo + a4 + a2$
f	$oo6 + a2$	$ooxo + a2 = xxo$
g	$a4 + a2$	$a2 + a2$
h	$oo6$	xxo
i	$a4$	xxo
j	$a2$	0

how L wins

rule	{game} or subgame	result
6	$\mathcal{A}', o\mathcal{A}', o\mathcal{O}o', o\mathcal{O}', \mathcal{I}$ empty	
a	$o13, o15, \dots$	$o11, o13 \dots$
b	$o11$	$o7 \text{ } xxo$
c	$o7$	$o5$
d	$o5$	xxo
7	$\mathcal{A}', o\mathcal{A}', o\mathcal{O}o', o\mathcal{O}', \mathcal{I}, \mathcal{O}'$ empty	
a	$oo8$	$ooxo \text{ } xxo$
b	xxo	0

how R can move



e.g. how R can move on oxoxoxoxo

can assume R clobbers to right

(oxoxoxoxo has left-right symmetry)

OXOXOXOXO $\rightarrow$ _OOXOXOXO

OXOXOXOXO $\rightarrow$ OX_OOXOXO

OXOXOXOXO $\rightarrow$ OXOX_OOXO

OXOXOXOXO $\rightarrow$ OXOXOX_OO

example game

a26 OX. . . OXOXOXOXOXOXOX

L $\rightarrow$ o23 OX. . . OXOXOXOXOXO XX

R $\rightarrow$ a12 oo10 OXOXOXOXOXOX OOXOXOXOXO

st.form a4 a2 oo10 OXOX OX OOXOXOXOXO

L $\rightarrow$ a4 a2 o7 OXOX OX OXOXOXO

R $\rightarrow$ o3 a2 o7 OOX OX OXOXOXO

L $\rightarrow$ a2 a2 o7 OX OX OXOXOXO

example game (cont.)

L $\rightarrow$ a2 a2 o7

ox ox oxoxoxo

st.form o7

oxoxoxo

R $\rightarrow$ a4

oxox oo

L $\rightarrow$ xxo

o xxo

R $\rightarrow$ a2

xo

L $\rightarrow$ 0

x



thank you

email hayward@ualberta.ca