# A Theory of Cloud Bandwidth Pricing for Video-on-Demand Providers

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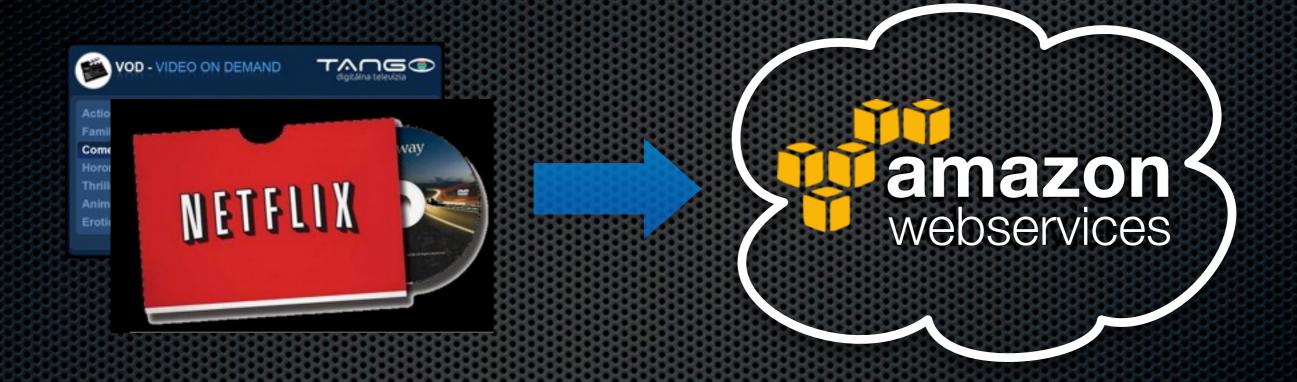
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#### A VoD Company Cares about

- Bandwidth guarantee
  - to maintain streaming quality
- Huge Bandwidth cost
  - "youku" pays > \$ 38.6M/year for bandwidth

# Can VoD leverage Cloud to reduce bandwidth cost with bandwidth guarantees

#### Current Cloud Platforms



Netflix moved to Amazon Web Services in 2010

Problem: No bandwidth guarantee!

## Extending Cloud-Tenant Interface

Bandwidth reservation for egress traffic of a VM

C. Guo et al.

SecondNet: a Data Center Network Virtualization Architecture with Bandwidth Guarantees ACM CONEXT '10

H. Ballani, P. Costa, T. Karagiannis, and A. Rowstron, Towards Predictable Datacenter Networks ACM SIGCOMM '11

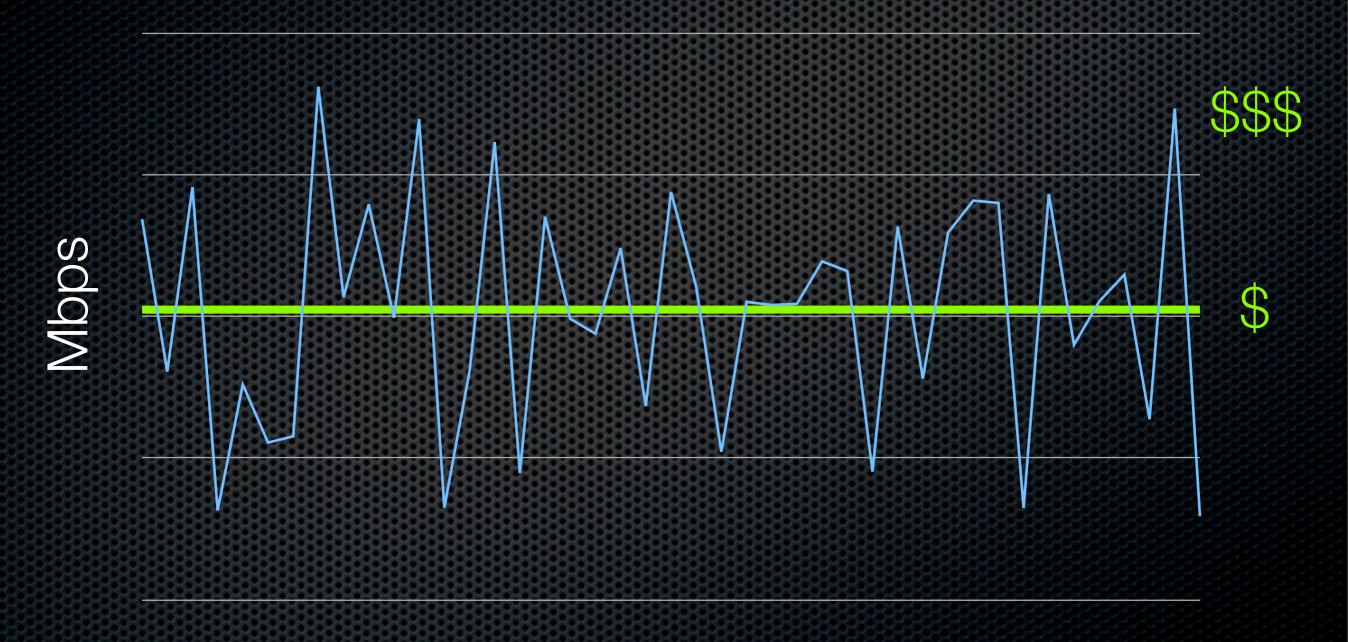
## How to Price Bandwidth Guarantees?

- VoD Companies want
  - bandwidth guarantee
  - reduced cost
- Cloud Providers want
  - economic incentives (profit)

### How about Pay-as-You-Go?

- Charge based on usage
  - number of bytes transferred in the past hour
- Not suitable for pricing bandwidth guarantees

## Example: Bursty vs. Constant



Time

#### Our Contributions

- A New Model
  - to price bandwidth guarantees
- Theoretical Analysis
  - to set pricing policy properly and fairly
- Practical Performance
  - simulations based on workload traces

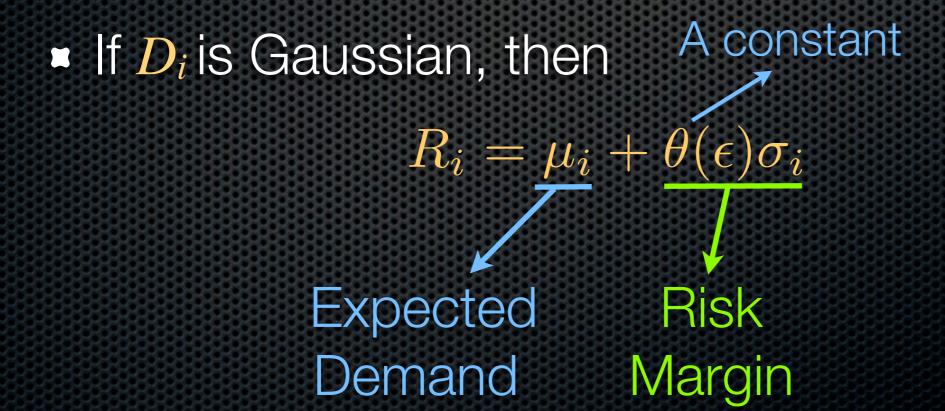
## Assumptions

- Each tenant i:
  - $\blacksquare$  random demand  $D_i$
  - $\blacksquare$  expectation  $\mu_i = E[D_i]$
  - standard deviation  $\sigma_{i^2} = \text{var}[D_i]$
- Demand covariance matrix  $\Sigma = [\sigma_{ij}]$

#### Probabilistic Bandwidth Guarantee

• With random demand  $D_i$ , tenant i wants to reserve bandwidth  $R_i$  such that

$$Pr(D_i > R_i) < \epsilon \longrightarrow Small constant$$



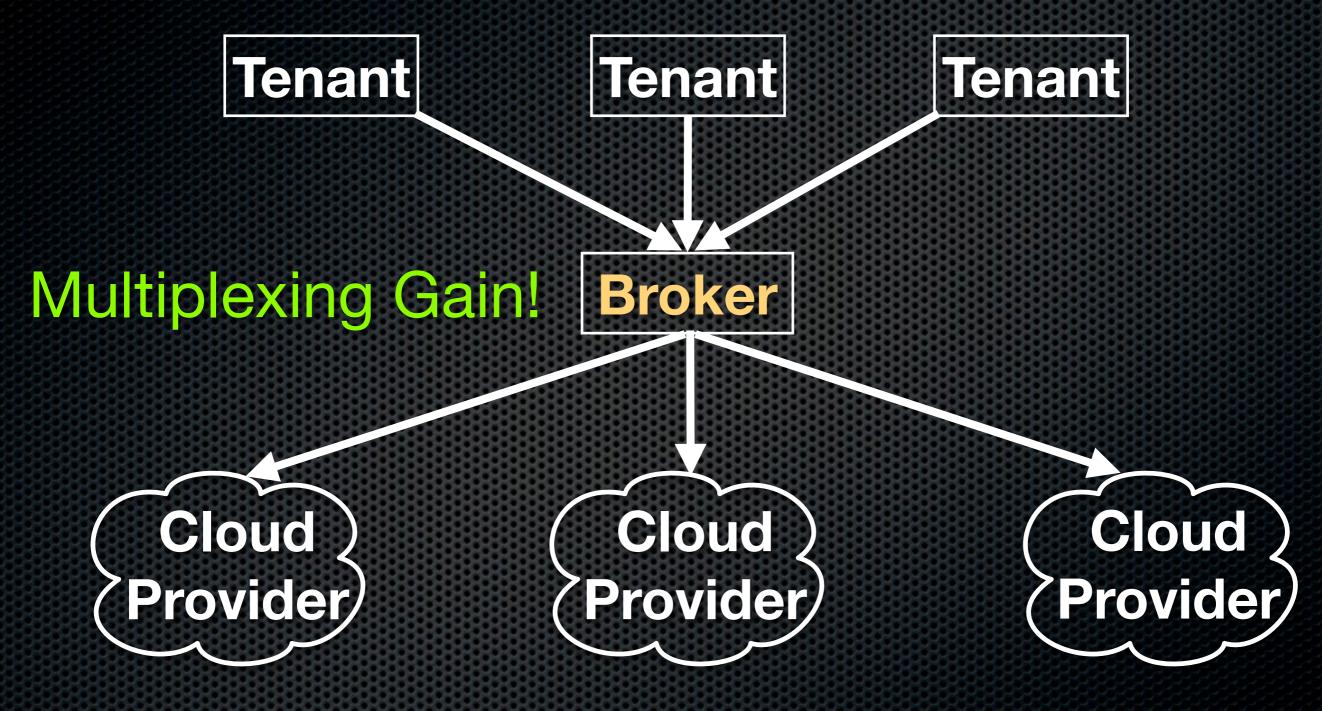
### Pricing Individual Reservation

- Assume: reserving 1 unit of bandwidth costs \$1
- Tenant i should pay

$$P_i = R_i = \mu_i + \theta(\epsilon)\sigma_i$$
Mean Burstiness Demand

- Individual reservation is costly
  - Most of the time, underutilization

#### An Economically Viable Broker



Reservation =  $\sum \mu_i$  + Shared Risk Margin Shared Risk Margin <  $\sum$  Individual Risk Margin

#### The Service Provided by the Broker

Individual Reservation  $R_i$ 

Probabilistic Guarantee  $\Pr(D_i > R_i) < \epsilon$ 

Random Demand

 $D_1$ 

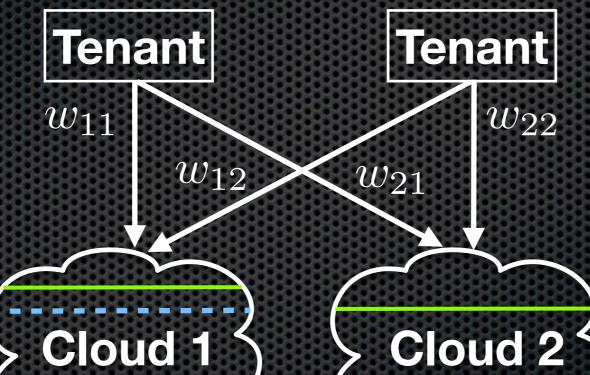
 $D_2$ 

Guaranteed Portion

 $w_i = \sum_s w_{si}$ 

Reservation  $A_1$ 

Load  $L_1 = \sum_i w_{1i} D_i$ 



Request Routing

 $A_2$  Reservation

Load  $L_2 = \sum_i w_{2i} D_i$ 

New Probabilistic Guarantee

$$\Pr(L_s > A_s) < \epsilon$$
, where  $L_s = \sum_i w_{si} D_i$ 

## Definition: Pricing Policy

- Guaranteed portion  $w_i = \sum_s w_{si}$
- Pricing strategy  $P_i(w_i, \mu_i, \sigma_i, \ldots)$ 
  - $\blacksquare P_i(w_i)$  increasing, concave,  $P_i = 0$  if  $w_i = 0$
- $\blacksquare$  Pricing policy  $\{P_i(\cdot)\}$
- Example: individual reservation

$$P_i = R_i = \mu_i + \theta(\epsilon)\sigma_i \longrightarrow P_i(w_i) = (\mu_i + \theta(\epsilon)\sigma_i)w_i$$

## Main Result 1 Pricing Region

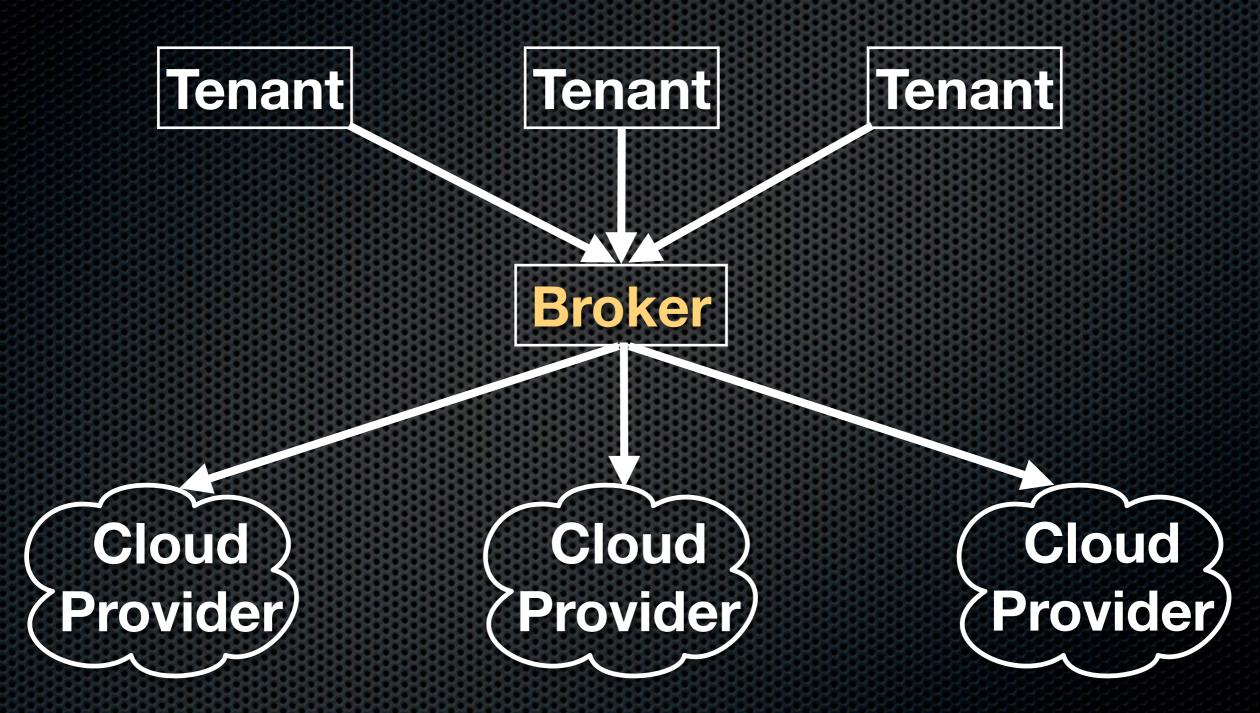
Objectives: Set prices such that tenants enjoy discount broker has incentive all demands guaranteed

## Good Pricing Region

- If and Only if  $P_i(w_i) \leq (\mu_i + \theta(\epsilon)\sigma_i)w_i$ 
  - Tenant i enjoys some price discount
- If and Only if  $P_i'(1) \ge \mu_i + \theta(\epsilon)\sigma_i\rho_{iM}$  to the market
  - A profit-driven broker is willing to guarantee all the tenant demand  $(w_i = 1, \forall i)$
- If both conditions are met:
  - Achieve Discount + Incentive + Guarantee!

 $\mu_i + \theta(\epsilon)\sigma_i\rho_{iM} \le P_i(1) \le \mu_i + \theta(\epsilon)\sigma_i$ 

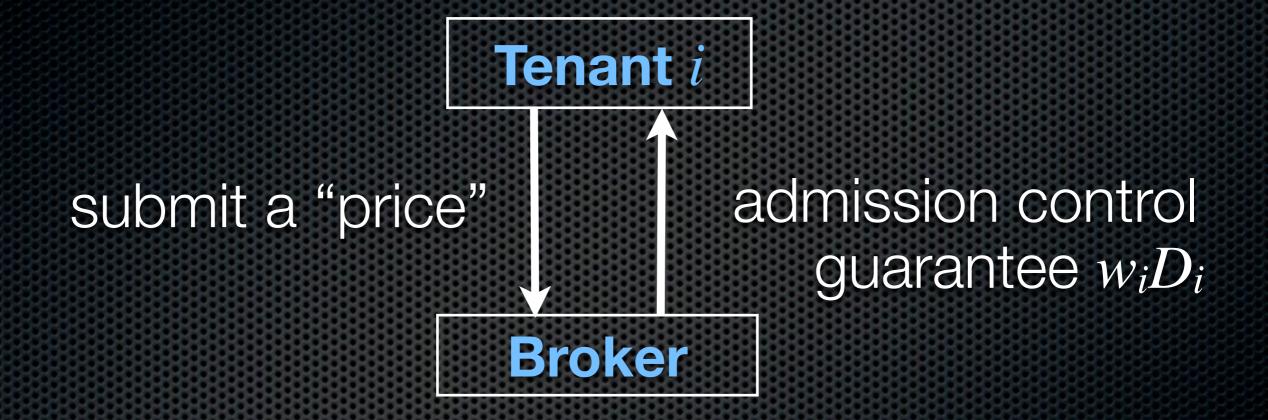
$$\mu_i + \theta(\epsilon)\sigma_i\rho_{iM} \le P_i(1) \le \mu_i + \theta(\epsilon)\sigma_i$$



With a good pricing policy, multiplexing gain is shared fairly.

How to enforce good pricing? Enforcement by an Authority \
What about a Free Market?

## Free Market: Allow Bargains



- \*  $w_i$  guaranteed portion in [0,1]
- **Submitted price is a function:**  $P_i(w_i)$
- If  $P_i(\cdot)$  is too low, the broker will set  $w_i < 1$ .

## A Game among all Tenants

$$\max \text{ Utility} = \left\{ \begin{array}{ll} -P_i(1), & \text{if } w_i = 1 \\ -\infty, & \text{if } w_i < 1 \end{array} \right.$$

Wants to reduce  $P_i(\cdot)$ , but must have  $w_i = 1$ 

#### Can prisoner's dilemma happen?

Pricing function  $P_i(\cdot)$ Strategy Players

Guarantee w<sub>i</sub>D<sub>i</sub>

#### **Broker**

Choose 
$$w_i$$
 to  $\max_{w_1,...,w_N} R = \sum_{i=1}^N P_i(w_i) - \sum_{s=1}^S A_s(w_1,...,w_N)$   
Profit Total Revenue Sum of Costs

## Main Result 2 Equilibrium Pricing

#### Equilibrium of the Free Market

With multiplexing, the market has a unique Nash equilibrium:

$$P_i^*(1) = \mu_i + \theta(\epsilon)\sigma_i\rho_{iM}$$
  
Expected Demand Correlation to the Demand Standard Deviation market, in [-1, 1]

Recall the good pricing region

$$\frac{\mu_i + \theta(\epsilon)\sigma_i\rho_{iM}}{I} \leq P_i(1) \leq \frac{\mu_i + \theta(\epsilon)\sigma_i}{I}$$
 Nash Equilibrium Individual Reservation

#### Trace-Driven Simulations

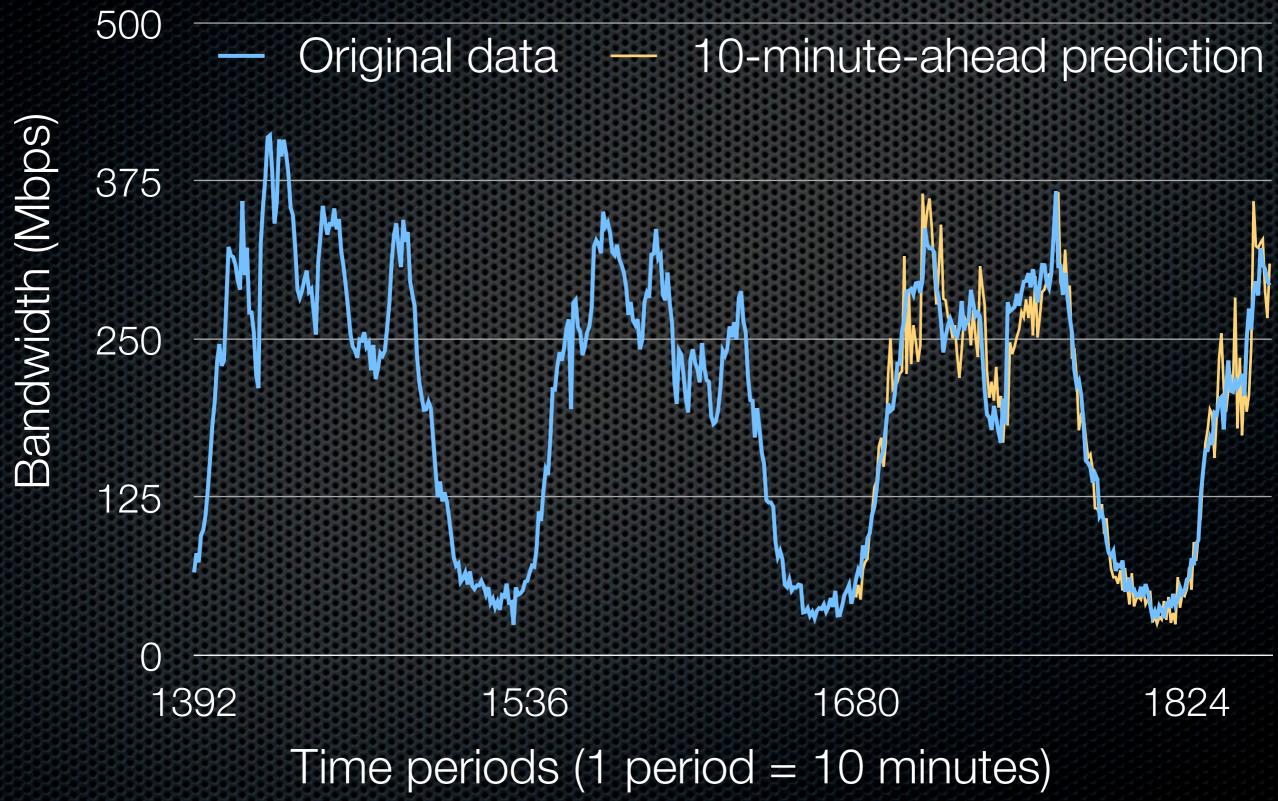
## Data Mining



- UUSee: a VoD provider in China
- Aggregate BW demand in video channels

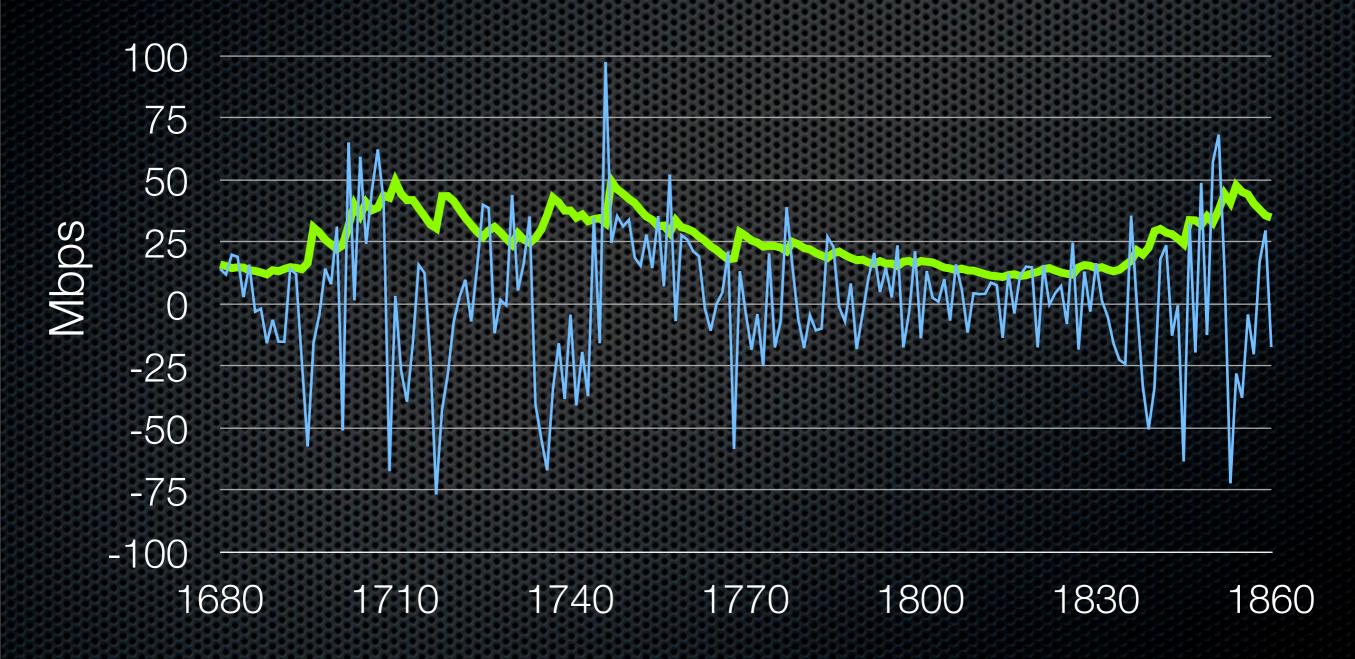
#### How to estimate expected demand?





#### How to estimate variance?

- Prediction error
- Estimated conditional error standard deviation



Time periods (1 period = 10 minutes)

#### Histogram of Price Discounts



Discounts of All Tenants in All Test Periods

#### Conclusions

- Broker Service:
  - Bandwidth reservation via multiplexing
- A theory to price bandwidth guarantees fairly and properly
- Pricing based on burstiness and correlation

## Thank you!

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