

# MILP formulation for open pit scheduling with multiple materials destinations

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## Abstract

*A mixed integer linear programming formulation for open pit production scheduling with multiple material destinations is presented.*

## 1. Notation

We will present an MILP formulation for the open pit production scheduling problem with multiple materials destinations. The notation of decision variables, parameters, sets, and constraints are as follows:

### 1.1. Sets

$\mathcal{K} = \{1, \dots, K\}$	set of all the mining-cuts in the model.
$\mathcal{P} = \{1, \dots, P\}$	set of all the phases (push-backs) in the model.
$\mathcal{D} = \{1, \dots, D\}$	set of all the possible destinations for materials in the model.
$\mathcal{A}_k = \{1, \dots, A_k\}$	set of all the directed arcs in the mining-cuts' precedence directed graph denoted by $G_k(\mathcal{K}, \mathcal{A}_k)$ .
$C_k(L)$	for each mining-cut $k$ , there is a set $C_k(L) \subset \mathcal{K}$ defining the immediate predecessor mining-cuts that must be extracted prior to extracting mining-cut $k$ . Where $L$ is an integer number presenting the total number of blocks in the set $C_k$ .
$B_p(M)$	for each phase $p$ there is a set $B_p(M) \subset \mathcal{K}$ defining the mining-cuts constructing phase $p$ . Phases are constructed to be sets of mining-cuts partitioning $\mathcal{K}$ . where $M$ is an integer number denoting the total number of blocks in the set $B_p$ .
$H_p(R)$	for each phase $p$ , there is a set $H_p(R) \subset \mathcal{K}$ defining the mining-cuts within the immediate predecessor pit phases (push-backs) that must be extracted prior to extracting phase $p$ . Where $R$ is an integer number presenting the total number of blocks in the set $H_p$ .

## 1.2. Indices

A general parameter can take the following indices in the format of  $f_{k,p}^{d,t,e}$ . Where:

$k \in \{1, \dots, K\}$	index for mining-cuts.
$p \in \{1, \dots, P\}$	index for phases.
$d \in \{1, \dots, D\}$	index for possible destinations for materials.
$t \in \{1, \dots, T\}$	index for scheduling periods.
$e \in \{1, \dots, E\}$	index for elements of interest in each block.

## 1.3. Parameters

$u_k^{d,t}$	the discounted dollar value generated by extracting mining-cut $k$ and sending it to destination $d$ in period $t$ .
$v_k^{d,t}$	the discounted revenue generated by selling the final products within mining-cut $k$ in period $t$ , if it is sent to destination $d$ , minus the extra discounted cost of mining all the material in mining-cut $k$ as ore and processing at destination $d$ .
$q_k^{d,t}$	the discounted cost of mining the material in mining-cut $k$ in period $t$ as waste and sending it to destination $d$ .
$g_k^e$	average grade of element $e$ in ore portion of mining-cut $k$ .
$\bar{g}^{-d,t,e}$	upper bound on acceptable average head grade of element $e$ , in period $t$ , at processing destination $d$ .
$\underline{g}^{d,t,e}$	lower bound on acceptable average head grade of element $e$ , in period $t$ , at processing destination $d$ .
$o_k$	ore tonnage in mining-cut $k$ .
$w_k$	waste tonnage in mining-cut $k$ .
$\bar{p}^{-d,t}$	upper bound on processing capacity of ore in period $t$ at destination $d$ (tonnes).
$\underline{p}^{d,t}$	lower bound on processing capacity of ore in period $t$ at destination $d$ (tonnes).
$\bar{m}^{-t}$	upper bound on mining capacity in period $t$ (tonnes).
$\underline{m}^t$	lower bound on mining capacity in period $t$ (tonnes).
$r^{d,e}$	processing recovery, is the proportion of element $e$ recovered if it is processed at destination $d$ .
$s^{t,e}$	price in present value terms obtainable per unit of product (element $e$ ).
$cs^{t,e}$	selling cost in present value terms per unit of product (element $e$ ).
$cp^{d,t,e}$	extra cost in present value terms per tonne of ore for mining and processing at destination $d$ .
$cm^{d,t}$	cost in present value terms of mining a tonne of waste in period $t$ sending it to destination $d$ .

$h_k$	bench number corresponding to mining-cut $k$ , benches are numbered from the top of the pit towards the bottom accordingly.
$\bar{h}_p$	maximum acceptable lead between phase $p$ and $p+1$ . Where the lead is the number of benches by which the mining of a specified phase must be ahead of the next one.
$\underline{h}_p$	minimum acceptable lead required between phase $p$ and $p+1$ .
$M$	the total number of mining-cuts in the set $B_p(M)$ .
$R$	the total number of mining-cuts in the set $H_p(R)$ .
$L$	the total number of mining-cuts in the set $C_k(L)$ .

#### 1.4. Decision Variables

$x_k^{d,t} \in [0,1]$	continuous variable, representing the portion of mining-cut $k$ sent to processing destination $d$ , in period $t$ .
$y_k^{d,t} \in [0,1]$	continuous variable, representing the portion of mining-cut $k$ mined in period $t$ , and sent to destination $d$ .
$b_k^t \in \{0,1\}$	binary integer variable controlling the precedence of extraction of mining-cuts. $b_k^t$ is equal to one if extraction of mining-cut $k$ has started by or in period $t$ , otherwise it is zero.
$z_p^t \in \{0,1\}$	binary integer variable controlling the precedence of mining phases. $z_p^t$ is equal to one if extraction of phase $p$ has started by or in period $t$ , otherwise it is zero.

## 2. Economic block value modeling

The objective function of the MILP formulation is to maximize the net present value of the mining operation. Hence, we need to define a clear concept of economic value based on the amount of ore within mining-cuts, which can be mined selectively. The profit from mining depends on the value of the mining-cut based on its processing destination and the costs incurred in mining and processing it. The cost of mining a cut is a function of its location, which characterizes how deep the mining-cut is located relative to the surface and how far it is relative to its final dump. The spatial factor can be applied as a mining cost adjustment factor for each mining-cut according to its location to the surface. The discounted profit from mining-cut  $k$  is equal to the discounted revenue generated by selling the final product contained in mining-cut  $k$  minus all the discounted costs involved in extracting mining-cut  $k$ , this is presented by Eq. (1).

$$u_k^{d,t} = \underbrace{\left[ \sum_{e=1}^E o_k \times g_k^e \times r^{d,e} \times (s^{t,e} - cs^{t,e}) \right]}_{\text{discounted revenues}} - \underbrace{\left[ \sum_{e=1}^E o_k \times cp^{d,t,e} \right]}_{\text{discounted costs}} - [(o_k + w_k) \times cm^{d,t}] \quad \forall d \in \{1, \dots, D\} \quad (1)$$

For simplification purposes we denote:

$$v_k^{d,t} = \sum_{e=1}^E o_k \times g_k^e \times r^{d,e} \times (s^{t,e} - cs^{t,e}) - \sum_{e=1}^E o_k \times cp^{d,t,e} \quad \forall d \in \{1, \dots, D\}, \quad k \in \{1, \dots, K\} \quad (2)$$

$$q_k^{d,t} = (o_k + w_k) \times cm^{d,t} \quad \forall d \in \{1, \dots, D\}, \quad k \in \{1, \dots, K\} \quad (3)$$

### 3. Model

Objective function:

$$\max \sum_{d=1}^D \sum_{t=1}^T \sum_{p=1}^P \left( \sum_{k \in B_p} (v_k^{d,t} x_k^{d,t} - q_k^{d,t} y_k^{d,t}) \right) \quad (4)$$

Subject to:

$$\underline{g}^{d,t,e} \leq \sum_{p=1}^P \left( \frac{\sum_{k \in B_p} g_k^e o_k}{\sum_{k \in B_p} o_k} x_k^{d,t} \right) \leq \bar{g}^{d,t,e} \quad \forall t \in \{1, \dots, T\}, \quad d \in \{1, \dots, D\}, \quad e \in \{1, \dots, E\} \quad (5)$$

$$\underline{p}^{d,t} \leq \sum_{p=1}^P \left( \sum_{k \in B_p} o_k x_k^{d,t} \right) \leq \bar{p}^{d,t} \quad \forall t \in \{1, \dots, T\}, \quad d \in \{1, \dots, D\} \quad (6)$$

$$\underline{m}^t \leq \sum_{p=1}^P \left( \sum_{k \in B_p} (o_k + w_k) y_k^{d,t} \right) \leq \bar{m}^t \quad \forall t \in \{1, \dots, T\}, \quad d \in \{1, \dots, D\} \quad (7)$$

$$\sum_{d=1}^D x_k^{d,t} \leq \sum_{d=1}^D y_k^{d,t} \quad \forall k \in \{1, \dots, K\}, \quad t \in \{1, \dots, T\} \quad (8)$$

$$b_k^t - \sum_{d=1}^D \sum_{i=1}^t y_s^{d,i} \leq 0 \quad \forall k \in \{1, \dots, K\}, \quad t \in \{1, \dots, T\}, \quad s \in C_k(L) \quad (9)$$

$$\sum_{d=1}^D \sum_{i=1}^t y_k^{d,i} - b_k^t \leq 0 \quad \forall k \in \{1, \dots, K\}, \quad t \in \{1, \dots, T\} \quad (10)$$

$$b_k^t - b_k^{t+1} \leq 0 \quad \forall k \in \{1, \dots, K\}, \quad t \in \{1, \dots, T-1\} \quad (11)$$

$$\underline{h}_p \geq h_l b_l^t - h_j b_j^t \geq \bar{h}_p \quad \forall p \in \{1, \dots, P\}, \quad t \in \{1, \dots, T\}, \quad l \in B_p, \quad j \in B_{p+1} \quad (12)$$

$$R.z_p^t - \sum_{r=1}^R \sum_{d=1}^D \sum_{i=1}^t y_r^{d,i} \leq 0 \quad \forall p \in \{1, \dots, P\}, \quad t \in \{1, \dots, T\}, \quad r \in H_p(R) \quad (13)$$

$$\sum_{m=1}^M \sum_{d=1}^D \sum_{i=1}^t y_m^{d,i} - M.z_p^t \leq 0 \quad \forall p \in \{1, \dots, P\}, \quad t \in \{1, \dots, T\}, \quad m \in B_p(M) \quad (14)$$

$$z_p^t - z_p^{t+1} \leq 0 \quad \forall p \in \{1, \dots, P\}, \quad t \in \{1, \dots, T-1\} \quad (15)$$