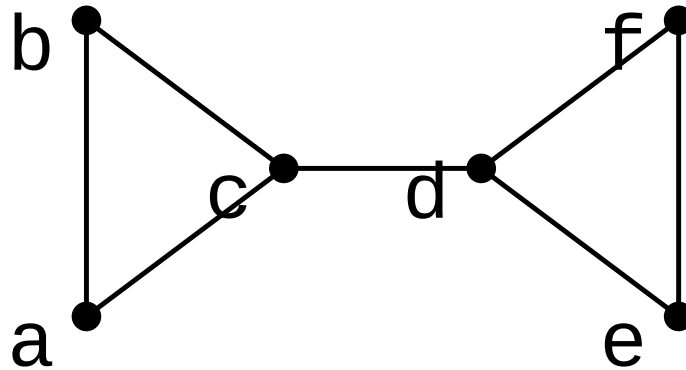
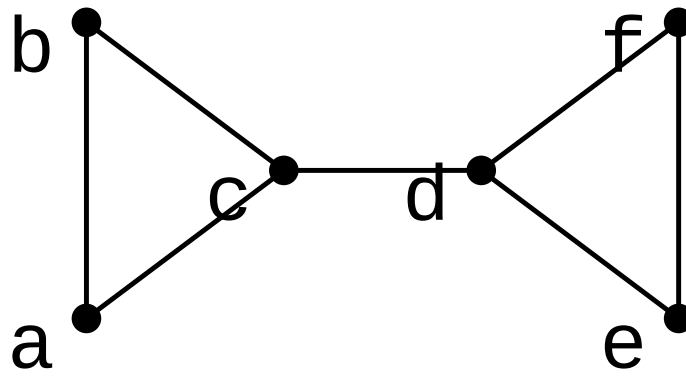


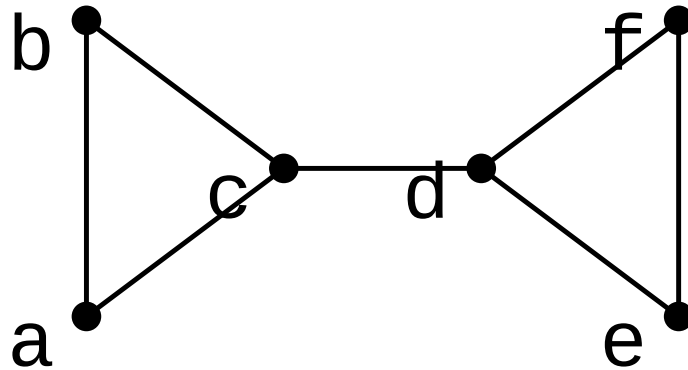


min cut

- * unweighted graph G , n nodes
- * cut: bipartition of node set
e.g. $\{a,d\} \{b,c,e,f\}$
- * an edge **crosses** a cut A,B
if it has one end in each part
e.g. $\{a,b\}$, $\{a,c\}$, $\{c,d\}$, $\{d,f\}$, $\{d,e\}$
- * cut size: number of edges that cross it
e.g. 5
- * min cut problem: given G , find any min cut



brute force min cut alg'm ?



brute force min cut alg'm

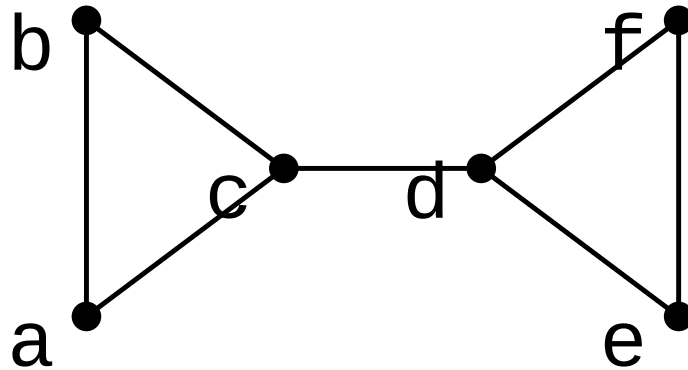
- try all possible cuts

e.g. {a} {bcdef}
 {ab} {cdef}
 {ac} {bdef}
 ...
 {abcde} {f}

total number of cuts, n nodes =
 (number subsets of n-1 set) - 1 = $2^{n-1} - 1$

[bipartition defined by making left part
 element a plus each partial subset]

here: $2^{n-1} - 1 = 2^5 - 1 = 31$



randomized kruskal min cut

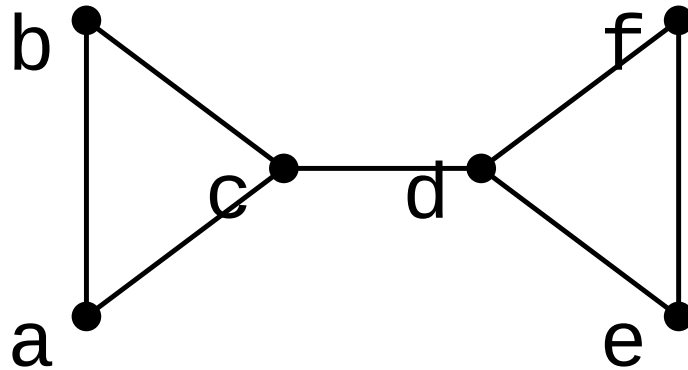
- * unweighted graph, n nodes
- * process edges in uniform-random order
- * apply Kruskal MST until exactly 2 components

edges (random order)	-	bc	ef	ac	ab	cd	de	ef
number components:		6	5	4	3	3	2	

cut bipartition {a b c d} {e f}

cut cross-edges {de, df}

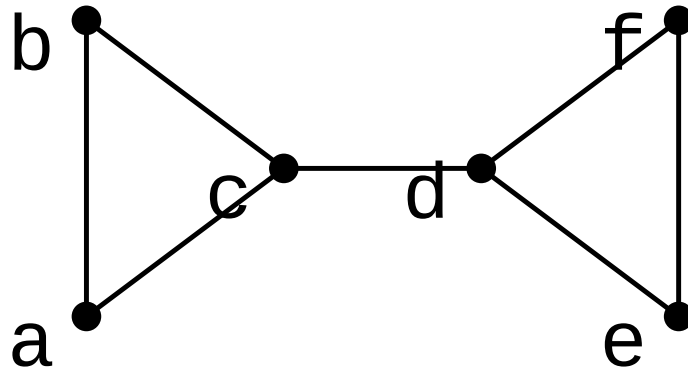
cut size 2



randomized kruskal min cut

- * unweighted graph, n nodes
- * consider edges in uniform-random order
- * apply Kruskal MST until exactly 2 components
- * $\text{prob}(\text{this cut } X \text{ is a min cut}) \geq 2/(n(n-1))$ (1)

proof of (1) ?



above example

* $\text{prob}(\text{RKMC cut is min}) \geq 2/(6 \cdot 5) = 1/15 = .0666\dots$

* $\text{prob}(\quad) = ?$

* unique min cut (UMC) $\{\{a, b, c\}, \{d, e, f\}\}$

* consider all $7!$ edge permutations

* $\text{RKMC}(\text{permutation})$ returns UMC iff

- in perm'n, edge cd is last, or

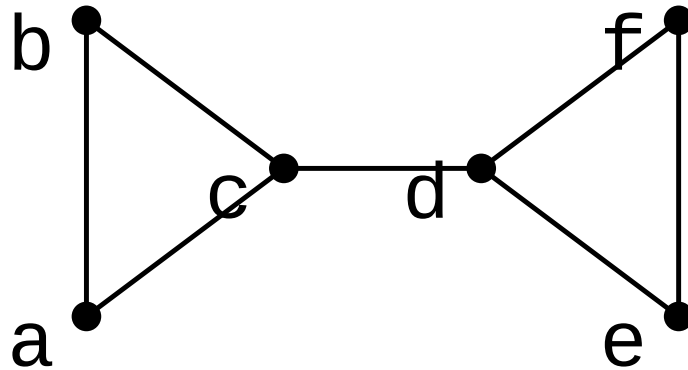
- in perm'n, edge cd is 2nd-last, or

- in perm'n, edge cd is 3rd-last and last

two edges are from different triangles

* $\text{prob} = 1/7 + 1/7 + (1/7)(3/5)$ exercise

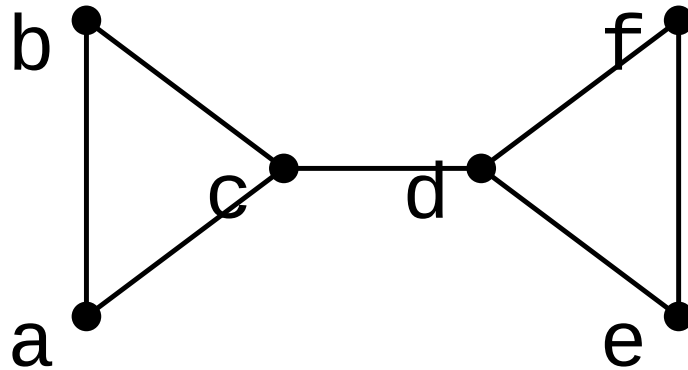
$= 13/35 = .3714\dots$



* call RKMC t times, take smallest cut found
 * prob(best cut found is min cut) ?

t	1.0 - (1.0 - 13/35)^t	
1	.371	
2	.604	with bound (1)
3	.751	67 trials
4	.843	prob >= .99
5	.901	
6	.938	with n = 100
7	.961	bound (1)
8	.975	22794 trials
9	.984	prob >= .99
10	.990	

after 10 trials
 prob(min cut) >= .99

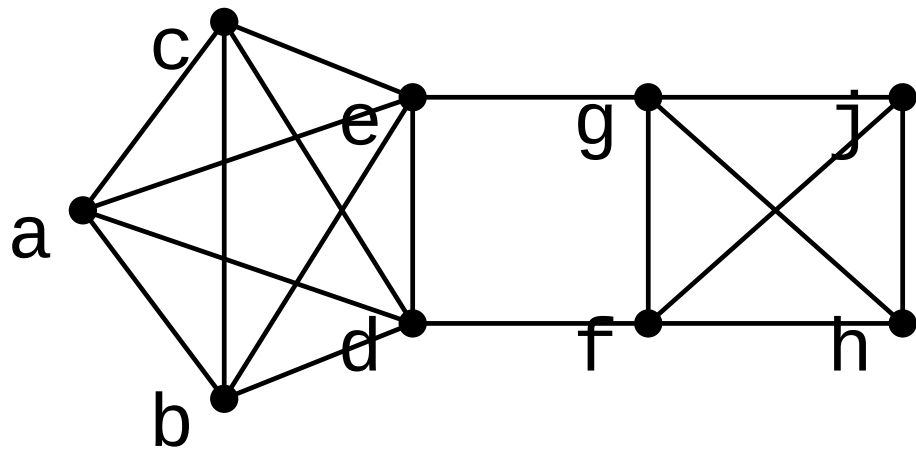


proof of (1)

- * F is RKMC forest-so-far
- * after t steps, F has $k = n-t$ components
- * number edges leaving component C_j of F ?
- * consider min cut with cross-edge set Y, $|Y| = y$
- * $\{C_j, V - C_j\}$ is a cut, so
 - at least y edges leave C_j
- * so number edges between components is $ky/2$
- * edges between components are exactly the RKMC-eligible edges
- * $\text{prob}(\text{next edge picked is in } Y)$

$$\leq y / (ky/2) = 2/k$$
- * $\text{prob}(\text{RKMC picks no edge from } Y)$

$$\geq \frac{n-2}{n} \frac{n-3}{n-1} \frac{n-4}{n-2} \dots \frac{3}{5} \frac{2}{4} \frac{1}{3} = \frac{2}{n(n-1)}$$



exercise

