

these notes are from

Chapter 7

Linear Programming and Reductions

in the class text

<http://algorithmics.lsi.upc.edu/docs/Dasgupta-Papadimitriou-Vazirani.pdf>

7.1.1 example: chocolate production

production * x1 boxes/day of type P profit 1/box
 * x2 " N 6

demand * P at most 200 boxes/day
 * N 300
 * P+N 400

maximize x1 + 6x2

subject to x1 <= 200
 x2 <= 300
 x1 + x2 <= 400
 x1, x2 >= 0
 x1, x2 rational

.....

vector/matrix notation

c		1		x		x1		A		1	0		b		200	
		6				x2				0	1				300	
										1	1				400	

maximize $c^T x$
such that $A x \leq b$
 $x \geq 0$
 x rational

7.1.1 example: chocolate production

production * x1 boxes/day of type P profit 1/box
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demand * P at most 200 boxes/day
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maximize $x_1 + 6x_2$

subject to $x_1 \leq 200$
 $x_2 \leq 300$
 $x_1 + x_2 \leq 400$
 $x_1, x_2 \geq 0$
 x_1, x_2 rational

solution: $x_1 = 100$, $x_2 = 300$, objective = 1900

proof of optimality?

$x_2 \leq 300$	$5x_2 \leq 5 \cdot 300 = 1500$
$x_1 + x_2 \leq 400$	$x_1 + x_2 \leq 400$

sum	$x_1 + 6x_2 \leq 1900$

magic LP trick :)

expanded problem: augment c, x, A, b

c		1		x		x1		A		1	0	0		b		200	
		6				x2				0	1	0				300	
		13				x3				1	1	1				400	
										0	1	3				600	

$$\begin{array}{ll}\max & c^T x \\ \text{s.t.} & A x \leq b \\ & x \geq 0\end{array}$$

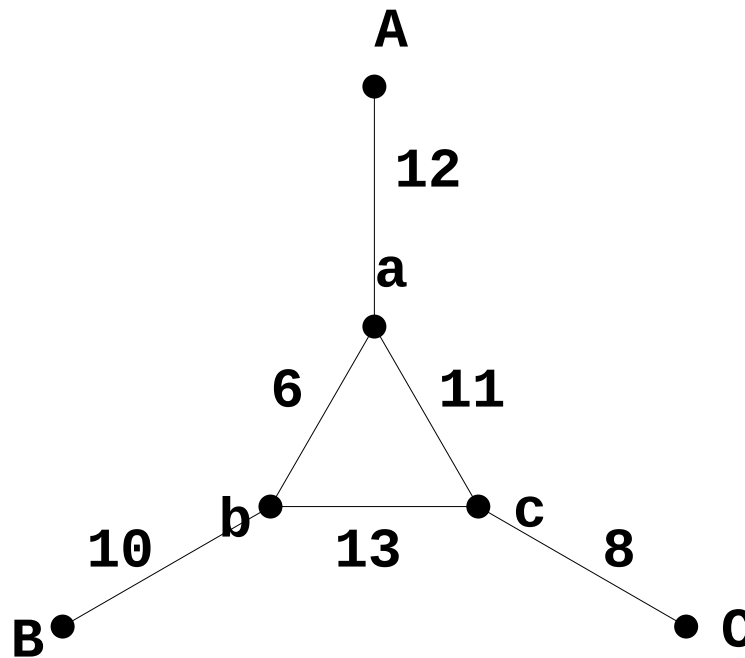
- evaluate at <https://sagecell.sagemath.org/>

```
p = MixedIntegerLinearProgram()
v = p.new_variable(real=True, nonnegative=True)
c, x, y, z = v["c"], v["x"], v["y"], v["z"]
p.set_objective(c)
p.add_constraint(c == x + 6*y + 13*z)
p.add_constraint(x <= 200)
p.add_constraint(y <= 300)
p.add_constraint(x + y + z <= 400)
p.add_constraint(y + 3*z <= 600)
round(p.solve(), 2)
p.get_values(c,x,y,z)
```

- you should get this output

```
[3100.0, 0.0, 300.0, 100.0]
```

7.1.3 example: optimum bandwidth allocation



- * users A B C
- * between each pair users, at least 2 units bandwidth
- * profit A-B \$3/unit bandwidth, B-C \$2/ub, A-C \$4/ub
- * inter-user communication routing:
 direct (3 hops) or indirect (4 hops)

$$\max 3x_{AB} + 3x'_{AB} + 2x_{BC} + 2x'_{BC} + 4x_{AC} + 4x'_{AC}$$

$$\begin{aligned} \text{s.t. } & x_{AB} + x'_{AB} + x_{BC} + x'_{BC} \leq 10 && \text{edge (b,B)} \\ & x_{AB} + x'_{AB} + x_{AC} + x'_{AC} \leq 12 && \text{edge (a,A)} \\ & x_{BC} + x'_{BC} + x_{AC} + x'_{AC} \leq 8 && \text{edge (c,C)} \\ & x_{AB} + x'_{BC} + x'_{AC} \leq 6 && \text{edge (a,b)} \\ & x'_{AB} + x_{BC} + x'_{AC} \leq 13 && \text{edge (b,c)} \\ & x'_{AB} + x'_{BC} + x_{AC} \leq 11 && \text{edge (a,c)} \\ & x_{AB} + x'_{AB} \geq 2 \\ & x_{BC} + x'_{BC} \geq 2 \\ & x_{AC} + x'_{AC} \geq 2 \\ & x_{AB}, x'_{AB}, x_{BC}, x'_{BC}, x_{AC}, x'_{AC} \geq 0 \end{aligned}$$

7.1.4 variants, standard inequality form

* max cx ... can be max or min
 * s.t. $3x_1 - 2x_9 \leq 3$ $\leq, \geq, =$
 * $x_1 \geq 0$ $\geq 0, \leq 0, \text{unrestricted}$

less-equal form:	standard form:	
max cx	max cx	
s.t. $Ax \leq b$	s.t. $Ax = b$	
$x \geq 0$	$x \geq 0$	

e.g. put this into
 standard less-equal form

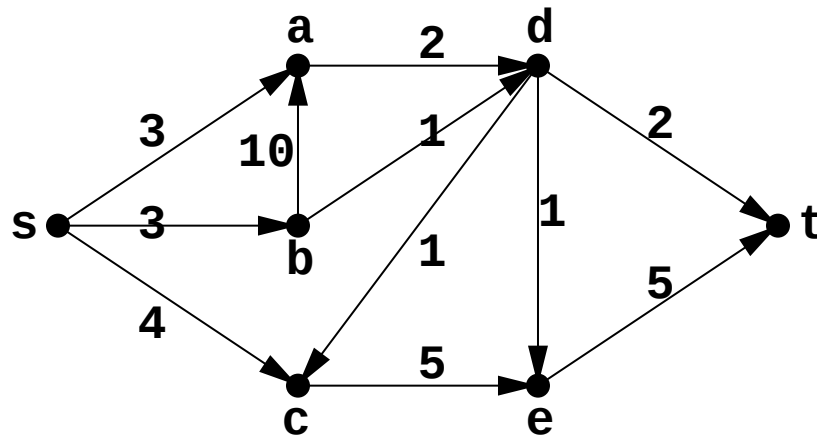
min $3x_1 - 4x_2$
 s.t. $x_1 + x_2 = 7$
 $x_1 - x_3 \geq 8$
 $x_1, x_3 \geq 0$

max $-3x_1 + 4x_2$
 s.t. $x_1 + x_2 \leq 7$
 $-x_1 - x_2 \leq -7$
 $-x_1 + x_3 \leq -8$
 $x_1, x_3 \geq 0$

max $-3x_1 + 4(x_4 - x_5)$
 s.t. $x_1 + (x_4 - x_5) \leq 7$
 $-x_1 - (x_4 - x_5) \leq -7$
 $-x_1 + x_3 \leq -8$
 $x_1, x_3, x_4, x_5 \geq 0$

max $-3x_1 + 4x_4 - 4x_5$
 s.t. $x_1 + x_4 - x_5 \leq 7$
 $-x_1 - x_4 + x_5 \leq -7$
 $-x_1 + x_3 \leq -8$
 $x_1, x_3, x_4, x_5 \geq 0$

7.2.2 max flow



network input

- * source s, terminus t
- * on each arc, capacity

flow

- * on each arc, non-negative flow $f(\text{arc})$ (e.g. 73 litre/s) that does not exceed $\text{capacity}(\text{arc})$
- * at each node except for source and terminus,
 - * $\text{in-flow}(\text{node})$ [sum, over in-arcs j, of $f(j)$] = $\text{out-flow}(\text{node})$ [sum, over out-arcs k, of $f(k)$]

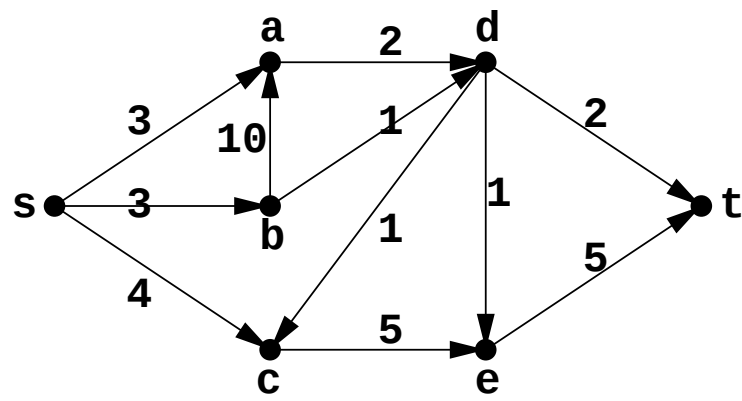
flow volume

- * $\text{out-flow}(\text{source})$, or $\text{in-flow}(\text{terminus})$

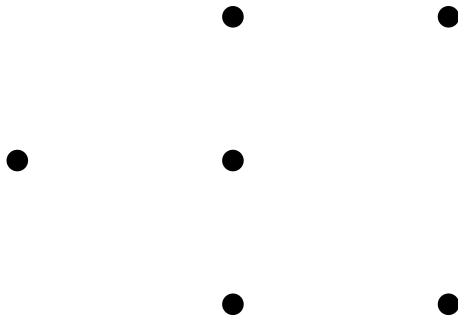
problem:

- * given network, find a max volume flow

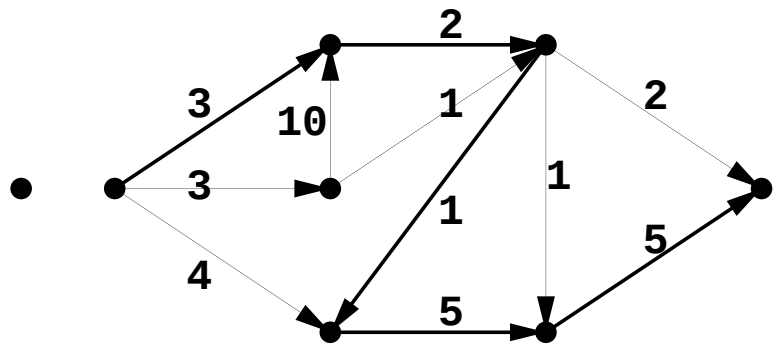
network



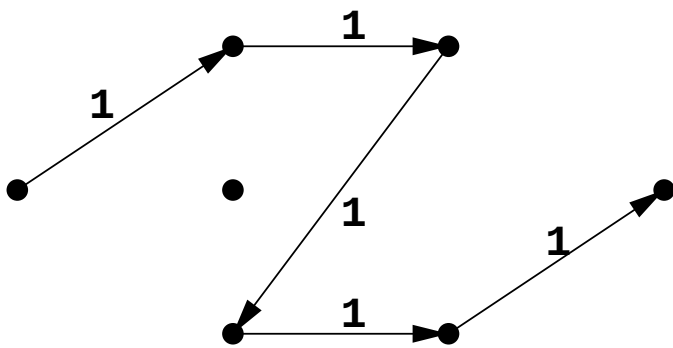
flow 0



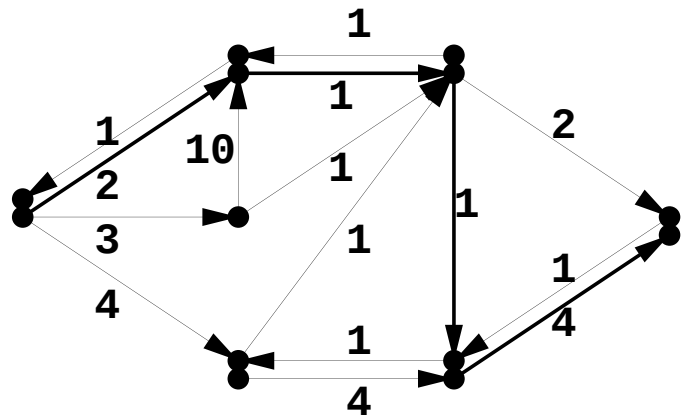
residual 0



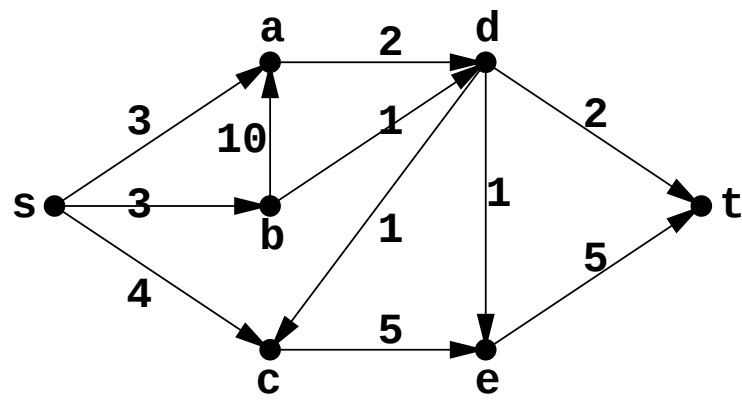
flow 1



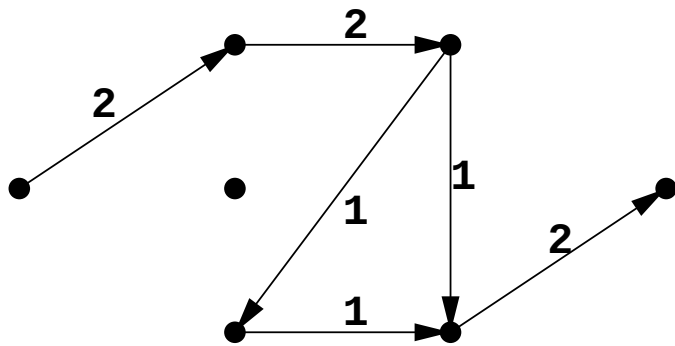
residual 1



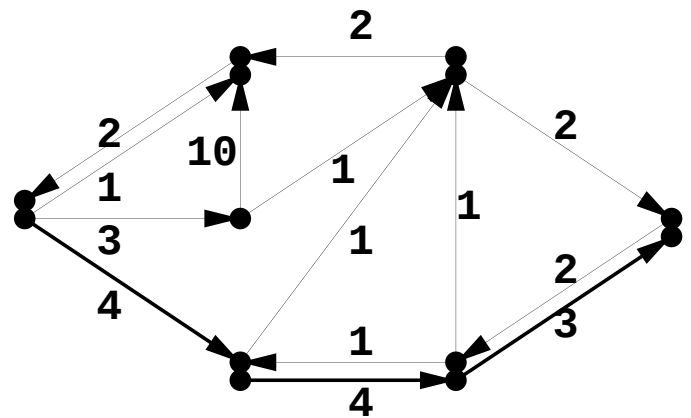
network



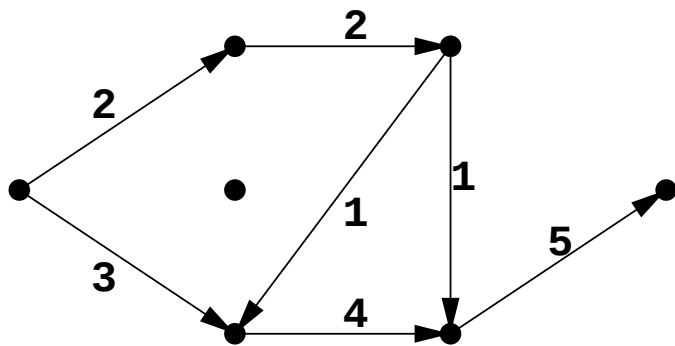
flow 2



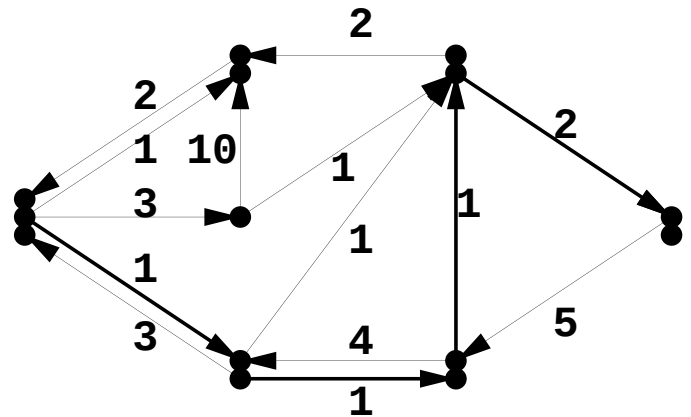
residual 2



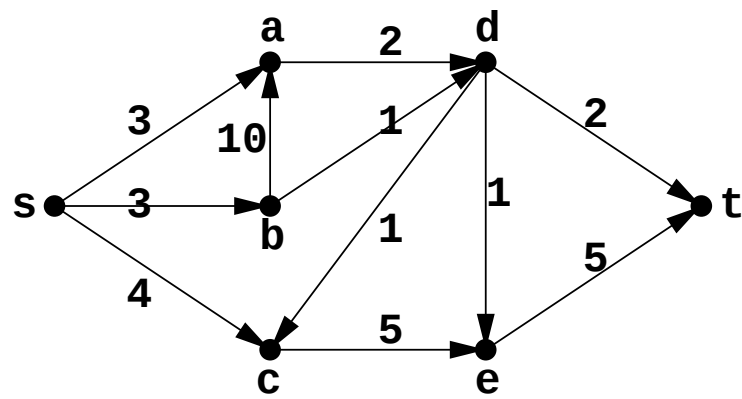
flow 3



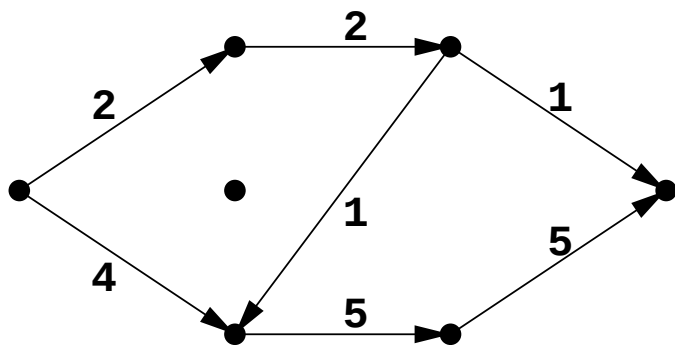
residual 3



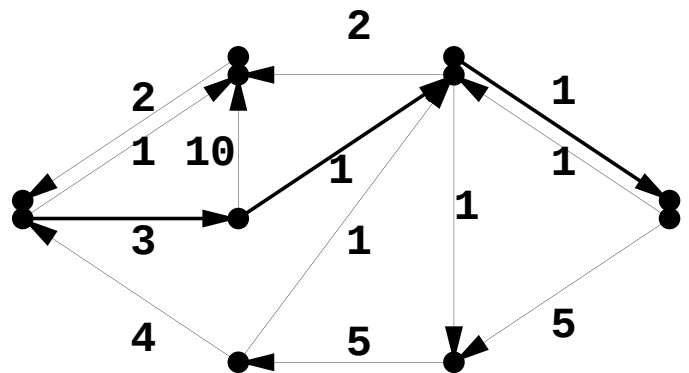
network



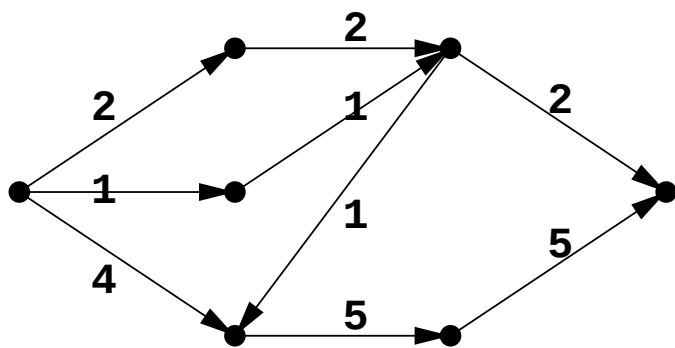
flow 4



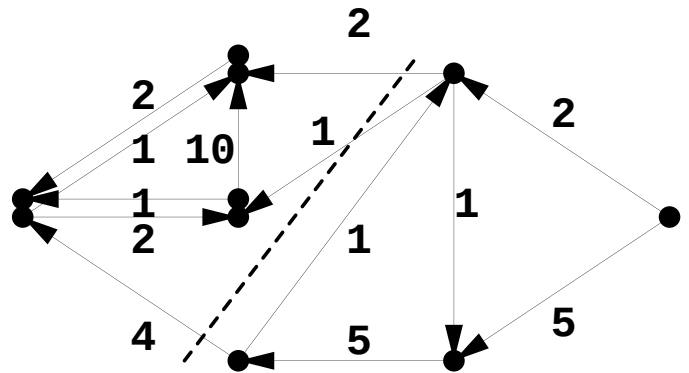
residual 4



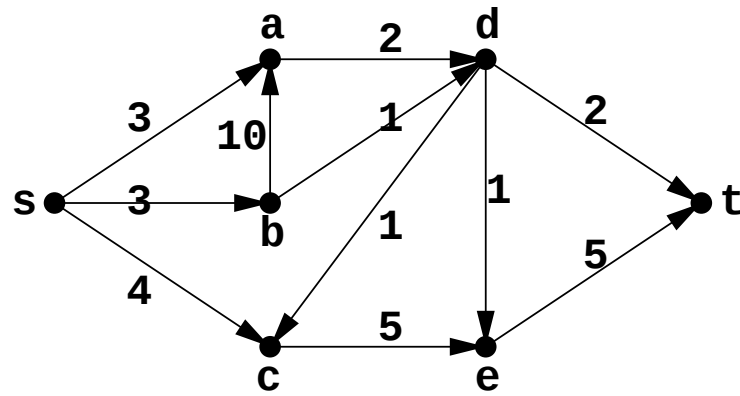
flow 5



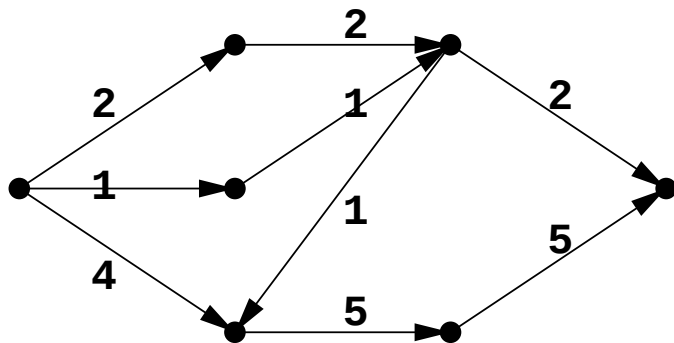
residual 5



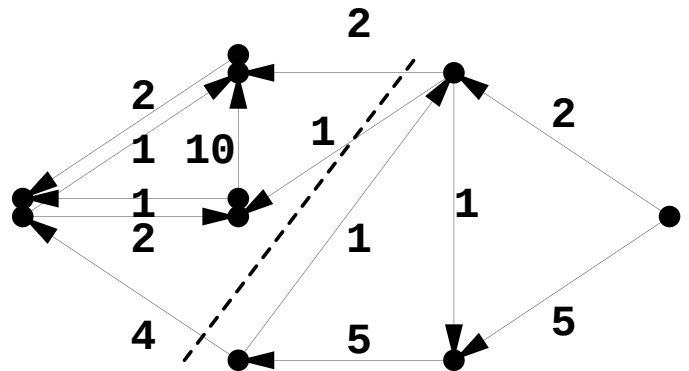
network



flow 5



residual 5



min cut?

in final residual network,

{ nodes reachable from s } { other nodes }

here

{ s, a, b } { c, d, e, t }

why?

7.2.2 max flow transforms to LP

variables sa sb sc ad ba bd ce dc de dt et

max $sa + sb + sc$ flow out of source

flow conservation

s.t. $sa + ba = ad$ node a

$sb = ba + bd$ b

$sc + dc = ce$ c

$ad + bd = dc + de + dt$ d

$ce + de = et$ e

capacities

$sa \leq 3$

$sb \leq 3$

$sc \leq c$

$ad \leq 2$

$ba \leq 10$

$bd \leq 1$

$ce \leq 5$

$dc \leq 1$

$de \leq 1$

$dt \leq 2$

$et \leq 5$

$sa, sb, sc, ad, ba, bd, ce, dc, de, dt, et \geq 0$

conclusion: max flow polytime transforms into LP

max flow integrality observation:

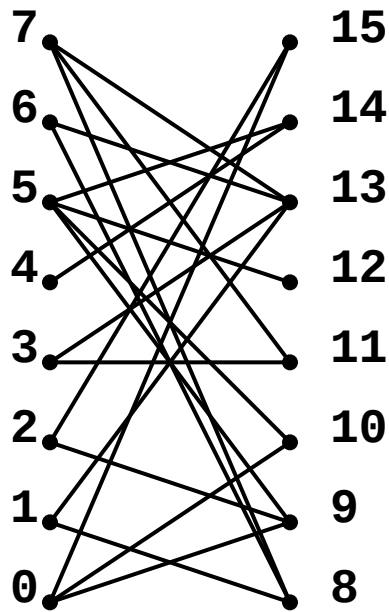
if all capacities are integer, then there exists

an all-integer max flow (why?)

7.3 max bipartite matching

instance: bipartite graph $G = (V_0, V_1, E)$

query: what is the size of a maximum matching in G ?



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