

family name _____ given name _____ ID _____

quiz1 2013oct4 c204 10marks 30min no electronic devices

1. [1.5 marks] For each $f(n)$, give the simplest $g(n)$ so that $f(n) = \Theta(g(n))$.

(i) $16n^2 \lg n + 11n^{5/2} + \frac{\sqrt{2}n^3}{(\lg n)^7}$ $g(n) =$ _____

(ii) $1.6n^2(\lg n)^5 + 5^{\lg n}$ $g(n) =$ _____

(iii) $39.7 \sum_{j=1}^n j^2 \lg j$ $g(n) =$ _____

2. [1.5 marks] Find some c so that, for all integers $n \geq 1$, $16n \lg n + 31n \leq cn^2$.

$c =$ _____

Justification:

3. [2 marks] Compute $2^{21} \pmod{13}$. Show your work.

so $2^{21} \pmod{13} =$ _____

```

def f(x,y): # y >= 0
    if y==0:
        return 0
    elif 0== y%2:
        return 2*f(x, y/2)
    else:
        return x + 2*f(x, y/2)

```

4. [2.5 marks] Claim. For all integers x , for all integers $y \geq 0$, $f(x,y)$ returns $x * y$.

Proof by induction on y . In the base case, $y = \underline{\hspace{2cm}}$, so in the execution of $f(x,y)$, the value returned is $\underline{\hspace{2cm}}$, which is equal to $x * y$, so the claim holds. Next consider the inductive case. Assume that the claim holds for all values of y in the set $\{0, \dots, t\}$. We want to show that the claim holds when $y = \underline{\hspace{2cm}}$. First assume y is even. Then the value returned is $\underline{\hspace{2cm}}$, which equals $\underline{\hspace{2cm}}$, which equals $\underline{\hspace{2cm}}$, since by the inductive hypothesis the claim holds when y has the value $\underline{\hspace{2cm}}$. So the claim holds in this case. [omit the rest of the proof]

5. [1 marks] Give a recurrence relation for the runtime $T(n)$ of $f(x,y)$, where n is the number of bits in x plus the number of bits in y .

$$\begin{aligned}
 T(n) &= \Theta(1) && \text{when } n = 0 \text{ (i.e. } y = 0) \\
 &= O(n) + T(\underline{\hspace{2cm}}) && \text{when } n \geq 1 \text{ and } y \text{ is even} \\
 &= O(\underline{\hspace{2cm}}) + T(\underline{\hspace{2cm}}) && \text{when } n \geq 1 \text{ and } y \text{ is odd}
 \end{aligned}$$

6. [1.5 marks] $f(x,y)$ is implemented efficiently. Multiplication of a k -bit number by 2 takes $O(\underline{\hspace{2cm}})$ time, and addition of two at-most- k -bit numbers takes $O(\underline{\hspace{2cm}})$ time. Assume that x and y each have $n/2$ bits. So, the total number of function calls resulting from $f(x,y)$ is $\Theta(\underline{\hspace{2cm}})$, each call takes $O(\underline{\hspace{2cm}})$ time, so the runtime of $f(x,y)$ is $O(\underline{\hspace{2cm}})$.

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(i) $16n \lg n + 11n^{3/2} + \frac{\sqrt{2}n^2}{(\lg n)^7}$ $g(n) =$ _____

(ii) $1.6n^2(\lg n)^5 + 7^{\lg n}$ $g(n) =$ _____

(iii) $39.7 \sum_{j=1}^n j \lg j$ $g(n) =$ _____

2. [1.5 marks] Find some c so that, for all integers $n \geq 1$, $12n \lg n + 26n \leq cn^2$.

$c =$ _____

Justification:

3. [2 marks] Compute $3^{21} \pmod{13}$. Show your work.

so $3^{21} \pmod{13} =$ _____

```

def f(x,y): # y >= 0
    if y==0:
        return 0
    elif 0== y%2:
        return 2*f(x, y/2)
    else:
        return x + 2*f(x, y/2)

```

4. [2.5 marks] Claim. For all integers x , for all integers $y \geq 0$, $f(x,y)$ returns $x * y$.

Proof by induction on y . In the base case, $y = \underline{\hspace{2cm}}$, so in the execution of $f(x,y)$, the value returned is $\underline{\hspace{2cm}}$, which is equal to $x * y$, so the claim holds. Next consider the inductive case. Assume that the claim holds for all values of y in the set $\{0, \dots, t\}$. We want to show that the claim holds when $y = \underline{\hspace{2cm}}$. First assume y is even. Then the value returned is $\underline{\hspace{2cm}}$, which equals $\underline{\hspace{2cm}}$, which equals $\underline{\hspace{2cm}}$, since by the inductive hypothesis the claim holds when y has the value $\underline{\hspace{2cm}}$. So the claim holds in this case. [omit the rest of the proof]

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$$\begin{aligned}
 T(n) &= \Theta(1) && \text{when } n = 0 \text{ (i.e. } y = 0) \\
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 \end{aligned}$$

6. [1.5 marks] $f(x,y)$ is implemented efficiently. Multiplication of a t -bit number by 2 takes $O(\underline{\hspace{2cm}})$ time, and addition of two at-most- t -bit numbers takes $O(\underline{\hspace{2cm}})$ time. Assume that x and y each have $n/2$ bits. So, the total number of function calls resulting from $f(x,y)$ is $\Theta(\underline{\hspace{2cm}})$, each call takes $O(\underline{\hspace{2cm}})$ time, so the runtime of $f(x,y)$ is $O(\underline{\hspace{2cm}})$.